

Please type a plus sign (+) inside this box → ☐

PTO/SB/05 (08-00)

Approved for use through 10/31/2002. OMB 0651-0032
U S Patent and Trademark Office: U S DEPARTMENT OF COMMERCE

Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number

UTILITY PATENT APPLICATION TRANSMITTAL

Attorney Docket No 19441 0002

First Inventor or Application Identifier Keng Tiong TAN

Title METHOD AND APPARATUS FOR NONLINEAR CODE-DIVISION
MULTIPLE-ACCESS TECHNOLOGY

(Only for new nonprovisional applications under 37 C.F.R. § 1.53(b))

Express Mail Label No.

APPLICATION ELEMENTS

See MPEP chapter 600 concerning design patent application contents.

1. ☒ Fee Transmittal Form (e.g., PTO/SB/17)
(Submit an original and a duplicate for fee processing)
2. ☐ Applicant claims small entity status.
See 37 CFR 1.27.
3. ☒ Specification [Total Pages 49]
(preferred arrangement set forth below)
 - Descriptive title of the Invention
 - Cross References to Related Applications
 - Statement Regarding Fed sponsored R & D
 - Background of the Invention
 - Brief Summary of the Invention
 - Brief Description of the Drawings (if filed)
 - Detailed Description
 - Claim(s)
 - Abstract of the Disclosure
4. ☒ Drawing(s) (35 U.S.C. 113) [Total Sheets 11]
5. Oath or Declaration [Total Pages]
 - a. ☐ Newly executed (original or copy)
 - b. ☐ Copy from a prior application (37 CFR 1.63 (d))
(for a continuation/divisional with Box 17 completed)☒ DELETION OF INVENTOR(S)
Signed statement attached deleting inventor(s)
named in the prior application, see 37 CFR
1.63(d)(2) and 1.33(b).
6. ☐ Application Data Sheet. See 37 CFR 1.76

ADDRESS TO

Assistant Commissioner for Patents
Box Patent Application
Washington, DC 20231

7. ☐ CD-ROM or CD-R in duplicate, large table or
Computer Program (Appendix)
8. Nucleotide and/or Amino Acid Sequence Submission
(if applicable, all necessary)
 - a. ☐ Computer Readable Form (CRF)
 - b. Specification Sequence Listing on:
 - i. ☐ CD-ROM or CD-R (2 copies); or
 - ii. ☐ paper
 - c. ☐ Statements verifying identity of above copies

ACCOMPANYING APPLICATIONS PARTS

9. ☐ Assignment Papers (cover sheet & document(s))
10. ☐ 37 C.F.R. § 3.73(b) Statement ☐ Power of
(when there is an assignee) Attorney
11. ☐ English Translation Document (if applicable)
12. ☐ Information Disclosure ☐ Copies of IDS
Statement (IDS)/PTO-1449 Citations
13. ☐ Preliminary Amendment
14. ☒ Return Receipt Postcard (MPEP 503)
(Should be specifically itemized)
15. ☐ Certified Copy of Priority Document(s)
(if foreign priority is claimed)
16. ☐ Other

17. If a CONTINUING APPLICATION, check appropriate box, and supply the requisite information below and in a preliminary amendment,
or in an Application Data Sheet under 37 CFR 1.76:

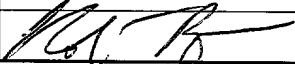
☐ Continuation ☐ Divisional ☐ Continuation-in-part (CIP) of prior application No _____ /
Prior application information: Examiner _____ Group / Art Unit _____

For CONTINUATION or DIVISIONAL APPS only: The entire disclosure of the prior application, from which an oath or declaration is supplied
under Box 5b, is considered a part of the disclosure of the accompanying or divisional application and is hereby incorporated by reference.
The incorporation can only be relied upon when a portion has been inadvertently omitted from the submitted application parts.

17. CORRESPONDENCE ADDRESS

☒ Customer Number or Bar Code Label 23517 or ☐ Correspondence address below
(Insert Customer No or Attach bar code label here)

Name	Edward Pennington				
	Swidler Berlin Shereff Friedman, LLP				
Address	3000 K Street, N.W.				
	Suite 300				
City	Washington	State	D.C.	Zip Code	20007-5116
Country	USA	Telephone	202 424-7500	Fax	202 424-7647

Name (Print/Type)	Robert C. Bertin	Registration No. (Attorney/Agent)	41,488
Signature		Date	November 15, 2000

Burden Hour Statement This form is estimated to take 0.2 hours to complete. Time will vary depending upon the needs of the individual case. Any
comments on the amount of time you are required to complete this form should be sent to the Chief Information Officer, U.S. Patent and Trademark
Office, Washington, DC 20231. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS SEND TO: Assistant Commissioner for
Patents, Box Patent Application, Washington, DC 20231.

JC903 U.S. PTO
09/712135

Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number

FEE TRANSMITTAL for FY 2001

Patent fees are subject to annual revision.

Complete if Known

Application Number
Filing Date November 15, 2000
First Named Inventor Keng Tiong TAN
Examiner Name
Group / Art Unit

TOTAL AMOUNT OF PAYMENT (\$ 1106

Attorney Docket No 19441.0002

METHOD OF PAYMENT (check one)

- 1.
- ☒
- The Commissioner is hereby authorized to charge indicated fees and credit any over payments to:

Deposit
Account
Number

195127, Order #19441.0002

Deposit
Account
Name

Swidler Berlin Shereff Friedman, LLP

- ☒ Charge Any Additional Fee Required
Under 37 CFR 1.16 and 1.17
☐ Applicant claims small entity status
See 37 CFR 1.27

- 2.
- ☐
- Payment Enclosed.

☐ Check ☐ Credit card ☐ Money
Order ☐ Other

FEE CALCULATION

1. BASIC FILING FEE

Large Fee Code	Entity Fee (\$)	Small Fee Code	Entity Fee (\$)	Fee Description	Fee Paid
101	710	201	355	Utility filing fee	710
106	320	206	160	Design filing fee	
107	490	207	245	Plant filing fee	
108	710	208	355	Reissue filing fee	
114	150	214	75	Provisional filing fee	

SUBTOTAL (1)

(\$ 710)

2. EXTRA CLAIM FEES

Total Claims	42	-20**	=	22	X	18	=	396
Independent Claims	2	-3**	=	0	X		=	0
Multiple Dependent					X		=	0

Large Fee Code	Entity Fee (\$)	Small Fee Code	Entity Fee (\$)	Fee Description
103	18	203	9	Claims in excess of 20
102	80	202	40	Independent claims in excess of 3
104	270	204	135	Multiple dependent claim, if not paid
109	80	209	40	** Reissue independent claims over original patent
110	18	210	9	** Reissue claims in excess of 20 and over original patent

SUBTOTAL (2)

(\$ 396)

**or number previously paid, if greater; For Reissues, see above

FEE CALCULATION (continued)

Large Fee Code	Entity Fee (\$)	Small Fee Code	Entity Fee (\$)	Fee Description	Fee Paid
105	130	205	65	Surcharge - late filing fee or oath	
127	50	227	25	Surcharge - late provisional filing fee or cover sheet	
139	130	139	130	Non-English specification	
147	2,520	147	2,520	For filing a request for reexamination	
112	920*	112	920*	Requesting publication of SIR prior to Examiner action	
113	1,840*	113	1,840*	Requesting publication of SIR after Examiner action	
115	110	215	55	Extension for reply within first month	
116	390	216	195	Extension for reply within second month	
117	890	217	445	Extension for reply within third month	
118	1,390	218	695	Extension for reply within fourth month	
128	1,890	228	945	Extension for reply within fifth month	
119	310	219	155	Notice of Appeal	
120	310	220	155	Filing a brief in support of an appeal	
121	270	221	135	Request for oral hearing	
138	1,510	138	1,510	Petition to institute a public use proceeding	
140	110	240	55	Petition to revive - unavoidable	
141	1,240	241	620	Petition to revive - unintentional	
142	1,240	242	620	Utility issue fee (or reissue)	
143	440	243	220	Design issue fee	
144	600	244	300	Plant issue fee	
122	130	122	130	Petitions to the Commissioner	
123	130	123	130	Petitions related to provisional applications	
126	180	126	180	Submission of Information Disclosure Stmt	
581	40	581	40	Recording each patent assignment per property (times number of properties)	
146	710	246	355	Filing a submission after final rejection (37 CFR § 1.129(a))	
149	710	249	355	For each additional invention to be examined (37 CFR § 1.129(b))	
179	710	279	355	Request for Continued Examination (RCE)	
169	900	169	900	Request for expedited examination of a design application	

Other fee (specify) _____


*Reduced by Basic Filing Fee Paid

SUBTOTAL (3)

(\$ 0)

SUBMITTED BY

Complete (if applicable)

Name (Print/Type) Robert C. Bertin Registration No Attorney/Agent 41,488 Telephone 202 424-7500
Signature  Date November 15, 2000

WARNING: Information on this form may become public. Credit card information should not be included on this form. Provide credit card information and authorization on PTO-2038.

Burden Hour Statement: This form is estimated to take 0.2 hours to complete. Time will vary depending upon the needs of the individual case. Any comments on the amount of time you are required to complete this form should be sent to the Chief Information Officer, U.S. Patent and Trademark Office, Washington, DC 20231. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS SEND TO: Assistant Commissioner for Patents, Washington, DC 20231.

**METHOD AND APPARATUS FOR NON-LINEAR CODE-DIVISION MULTIPLE
ACCESS TECHNOLOGY**

5

FIELD OF THE INVENTION:

The present invention relates generally to code division multiple access (CDMA) communications technology and, more particularly, to the selection and use of a class of non-linear spreading codes to improve performance characteristics of CDMA technology.

10

BACKGROUND OF THE INVENTION:

Recent advances in technology have given rise to communications electronics that are faster, consume less power and are less expensive as compared to earlier generations. This in turn has caused rapid growth in the global communications market, which includes both fixed and mobile segments. This rapid growth has manifested itself through increasing numbers of users of communications technologies, and the increasing services and bandwidth available to users. This growth is expected to continue for many years to come.

15

Current technologies for multi-user communication systems include code division multiple access (CDMA) and time division multiple access (TDMA), both of which are widely implemented in mobile communications. TDMA is used in the United States (IS-136) and Europe (GSM) as a digital wireless technology. CDMA (IS-95) has been implemented in digital wireless systems in the past few years and exhibits certain improved performance characteristics. CDMA accordingly appears poised to overtake TDMA and become the preferred technology for third generation mobile communications systems, which seek to provide high-speed data services in addition to providing high-quality voice services.

20

25

CDMA and TDMA have different performance characteristics in several areas. CDMA works by coding message bits into code sequences, which in turn are modulated for transmission over a wireless channel. In contrast to TDMA, the coding allows correction of some transmission errors due to noise in the channel, at least when there is less than full occupancy of the communication channel.

In a widely used mobile cellular implementation of CDMA, up to 64 (or 256) signals are transmitted in parallel from a base station to mobile units. In realistic noise environments, this number is limited by the peak power that can be transmitted by law or other considerations. There is accordingly a necessary balance between the transmitted signal power of the composite CDMA signal and the number of parallel CDMA active users supported. Although a higher transmitted signal power will usually result in a better coverage and signal reception at the receivers, this will also result in higher noise in neighboring cells. A performance indicator for mobile communication systems is the peak-to-average power (PAP) magnitude of the composite CDMA signals. High PAP has always been an inherent problem of CDMA systems. Pulse shaping and complex modulation techniques such as continuous phase modulation techniques have been developed to alleviate negative effects of high PAP. However, the problem persists.

Accordingly, there is a need for a new system and method for leveraging the advantages of CMDA that can increase the performance of CDMA to allow its operation in a third generation environment of high-data rates. There is a further need for a system which alleviates problems associated with high PAP. There is still a further need for such improvements in CDMA performance to be capable of low-cost implementation in hardware, software or firmware within existing CDMA systems to realize performance improvements.

SUMMARY OF THE INVENTION:

According to the present invention, the performance of code division multiple access technology is enhanced through the use of a new class of non-linear block codes during code division multiple access signal coding. This class of non-linear block codes is termed Go-CDMA codes and is defined in terms of its mathematical properties in the Detailed Description section. By way of summary, the Go-CDMA codes are $n \times \ell$ nonlinear block codes constructed using column-reduced and row-reduced Hadamard orthogonal matrices, termed Go-CDMA matrices. The parameters n, ℓ are positive integers: n represents chips of user data transmitted in frames of size $\ell \leq \alpha n$, where α is the frame expansion factor. The codes map n -vectors containing binary message data to binary or multi-level ℓ -vectors for transmission, where $\ell \geq n$.

Go-CDMA codes have application in code-division, multiple access communication systems, giving improved performance on many measures over conventional CDMA and TDMA systems. These measures include, for example, peak-to-average power ratio (PAP), error correction as a function of ℓ/n and n/A , channel capacity C in terms of message data rates, transmitted bit error rate as a function of signal-to-noise ratios (SNR), signal-to-interference ratio (SIR), where the interference is from neighboring cells, upper limits to the active-user numbers in a communication cell, and computational effort in coding and decoding.

According to an embodiment of the present invention, a method for coding a code division multiple access signal based on Go-CDMA codes, includes providing majority logic coding blocks, where each block comprises a Go-CDMA matrix. The method further includes coding a data message based on the majority coding blocks and transmitting the coded data message over a communication channel.

The majority coding blocks may comprise a single coding stage, two coding stages, three coding stages or more than three coding stages. The method may further comprise coding a plurality of data messages based on the majority coding blocks. The data messages may include at least one data message associated with an active user, at least one data message associated with a pseudo active user and/or at least one of the data messages associated with an inactive user. There may also be a permutation stage between each adjacent pair of coding stages depending on whether or not multiple coding stages are implemented. The majority coding logic blocks may also be implemented as a look up table. In this scenario, the coding is performed based on the look up table.

The data messages may include data elements in ternary format or in polar binary format. Moreover, each of the data messages may be derived from data received from an intermittent data source.

According to another embodiment of the present invention, a method for decoding a code division multiple access signal includes providing majority logic decoding blocks, where each block comprises a Go-CDMA matrix based on Go-CDMA codes. The method further includes receiving a signal over a communication channel and decoding a data message from the signal based on the majority coding blocks. This method is essentially the reverse of the coding process described above.

BRIEF DESCRIPTION OF THE FIGURES:

The above described features and advantages of the present invention will be more fully appreciated with reference to the detailed description and appended figures, in which:

Figure 1 depicts a communication channel with additive noise.

Figure 2 depicts a multiple access coding-decoding communication system according to an embodiment of the present invention.

Figures 3A and 3B depict functional block diagrams of a CDMA system used in mobile communications which incorporates coding and decoding blocks according to an embodiment of the present invention.

Figure 4 depicts an illustration of transmitted peak to average power for TDMA, CDMA and Majority Logic systems.

Figure 5 depicts a two stage Majority Logic coding scheme for encoding (or decoding with arrows reversed) up to nine active communications channels using 3×3 single stage Majority Logic codes according to an embodiment of the present invention.

Figure 6 depicts Majority Logic Coding using a Hadamard sub-matrix M according to an embodiment of the present invention.

Figure 7 depicts a Majority Logic decoding scheme for decoding of the i^{th} message received from the i^{th} transmitter according to an embodiment of the present invention.

Figure 8 depicts a two stage Go-CDMA coding (or decoding with arrows reversed) scheme for coding (or decoding) up to 25 active users using 16×5 single-stage Go-CDMA codes according to an embodiment of the present invention.

Figure 9 depicts a three stage Go-CDMA coding scheme for coding (or decoding with arrows reversed) up to 125 active users using 16×5 single-stage Go-CDMA codes according to an embodiment of the present invention.

Figure 10 depicts a four stage Go-CDMA coding scheme for coding (or decoding with arrows reversed) up to 625 active users using 16×5 single-stage Go-CDMA codes according to an embodiment of the present invention.

Figure 11 depicts active user data and psuedo-active user data as a data vector d according to an embodiment of the present invention.

Figure 12 depicts a method of coding data for transmission within a CDMA system according to an embodiment of the present invention.

5 Figure 13 depicts a method of decoding data received within a CDMA system according to an embodiment of the present invention.

Figure 14 depicts a method of generating Go-CDMA codes according to an embodiment of the present invention.

DETAILED DESCRIPTION:

The performance of code division multiple access (CDMA) technology is enhanced through the use of a new class of non-linear block codes during CDMA signal coding (and decoding). This class of non-linear block codes is termed Go-CDMA codes and is defined in terms of its mathematical properties in the Go-CDMA coding overview and Go-CDMA Matrices sections below. Prior to describing Go-CDMA coding and decoding according to the present invention, an overview of pertinent communications technologies is presented including an overview of those technologies in which Go-CDMA coding and decoding may be implemented.

I. Overview of Relevant Communication Technologies and Coding Schemes

Current technologies for single cell, or multiple cell, multi-user communication systems include CDMA and time-division-multiple-access (TDMA). These technologies are widely used for mobile communication, with TDMA being the basis of the GSM mobile telephone system used in Europe.

To illustrate such communications systems, Figure 1 illustrates an environment in which multi-user communication systems exist. Referring to Figure 1, a communication channel 100 is illustrated as having additive noise on it. The communication channel may be, for example, air, space, an electrical connection such as a wire, transmission line, or microwave element or an optical fiber. An incident signal s traversing the communications channel 100 is influenced by noise in the communications channel resulting in a signal $s + noise$ at the far end of the transmission line.

Figure 2 depicts a schematic of a multiple-access, coding-decoding communication system 200. The system 200 includes a multiple messages for transmission as inputs to a coding block 210. The coding block 210 encodes the messages and transmits the encoded messages as a

composite signal over the noisy communications channel 100. The decoding block 220 receives the composite signal which includes noise and decodes the encoded messages through a process that is in general the reverse of the coding process.

Both the coding block 210 and the decoding block 220 are depicted as single blocks. In the case of multiplexed optical fiber communications systems, for example, there may indeed be single coding and decoding blocks which interface with the optical fibre which is the communications channel 100. Alternatively, in mobile communications systems, for example, one or both of the coding block 210 and the decoding block 220 actually be implemented as multiple respective coding or decoding blocks, each with a unique spatial position relative to each other. This scenario is depicted in Figures 3A and 3B.

Referring to Figure 3A, a communications device is illustrated. The communications device may be any communications device including, for example, a base station 300 or a mobile communications unit 310 used in cellular communications. The device 300, 310 may include a modulation /demodulation unit 320 coupled to an antenna 340, coding and decoding units and 210 and 220 respectively, an optional pre and post coding and decoding unit 330, a processor 350, a memory 360 and I/O units 370.

The processor 350 may be a microprocessor, a micro-controller, a digital signal processor, an application specific integrated circuit or any other device suitable for controlling the operation of the device 300, 310. The processor 350 controls the operation of the device 300 and 310 and may be coupled to each of the functional blocks within the device to control their operation. Alternatively, each of the functional blocks depicted as within the device 300, 310 may not be implemented on devices separate from the processor. Rather, they may be functions performed

by the processor. The processor may control the device 300, 310 by executing program instructions stored in the memory 360.

The memory 360 stores data and may store program instructions for execution by the processor 350 or other elements within the device 300, 310. The memory may include volatile memory, non-volatile memory or both. The memory may include, for example read only memory (ROM) and read only memory devices such as CD-ROM devices, hard and floppy disk drives, random access memory (RAM), databases and any other type of memory or memory device.

The I/O units 370 may include any type of input output devices, including a display, a keyboard, a microphone, a speaker, a vibrating device, a modem for connecting to a network such as the PSTN, a local or wide area network or the interconnected network of servers, routers and bridges collectively known as the Internet.

During operation, the processor 350 may cause the device 300, 310 to open a communications channel via the antenna 340 with another communications device pursuant to the CDMA or TDMA protocol. In the case of wireless cellular telephony, the communications channel may be used to place a telephone call. The processor 350 also may receive signals from the I/O units 370 or the memory 360, such as voice or data signals, and may output data messages to the pre and post coding and decoding unit 330 based on the received data or voice signals. The data messages may in turn be sent through the coding unit 210 and the modulation/demodulation unit 320 and out the antenna 340 pursuant to the appropriate communications protocol. Similarly in the reverse direction, the processor may receive data messages via the antenna 340, the decoding unit 220, the pre and post decoding unit 330. The processor may then output a signal or other data, based on the received data messages, to one or

more of the I/O units 370 or may store the data in the memory 360. In this manner the device 300, 310 may perform communications functionality on behalf of a user of the device.

The pre and post coding and decoding block 330 is optional and may be used to, for example, insert (or decipher in the case of decoding) error correcting codes into the data messages, to interleave or de-interleave data or to otherwise manipulate the data messages prior to coding or after decoding. In the case of inserting error correction codes, any conventional error correction or error protection schemes may be used including cyclical redundancy check (CRC) schemes and forward error correction (FEC) schemes.

The coding and decoding blocks may be conventional CDMA or TDMA coding blocks. Alternatively, the coding and decoding blocks 210 and 220 may implement the Go-CDMA spreading code scheme for enhanced CDMA performance according to the present invention.

The modulation/demodulation unit 320 may be implemented with any appropriate amplifier to create a modulated output signal s based on either the TDMA or CDMA scheme, including CDMA schemes with Go-CDMA technology.

When a CDMA capable communications device 300, 310 includes a processor or other device that executes program instructions to perform the coding and decoding functions of blocks 210 and 220, the memory 360 may be updated with data and programming instructions to configure the coding and decoding blocks 210 and 220 to implement the Go-CDMA coding and decoding scheme according to the present invention. The program instructions and data may be loaded into the memory 360 via one or more of the I/O units 370 or via data received from the antenna 340.

Figure 3B depicts an illustrative view of a plurality of mobile units 310 engaged in cellular communications over a noisy air channel 100 with base stations 300. The mobile stations 310

and each of their respective coding units 210 may collectively be considered equivalent to the single coding unit 210 depicted in Figure 2 for coding n data messages for transmission over a noisy channel 100. In this scenario, the base station unit and its decoding unit 220 may be considered equivalent to the single decoder 220 depicted in Figure 2 for decoding n received data messages in a composite signal plus noise.

Several coding schemes for multiplexing data messages are conventionally used. Included among them are:

- Orthogonal-CDMA using linear Hadamard matrix block codes.
- Direct sequence, *pseudo*-random CDMA using linear *pseudo*-random codes.
- TDMA uses a (trivial) unity element of the class of orthogonal CDMA codes.

In the third generation standards for wideband mobile communication, a CDMA approach has been chosen. A latent technology is Majority Logic Coding which has not yet delivered significantly for any widely used communications system. These schemes are now discussed in terms of their properties.

System Capacity: For a given communication channel subject to noise, there is a theoretical upper data transmission rate, termed the Shannon-Hartley Channel Capacity, at which data can be transmitted error free. Practical schemes fall short of this limit. In a multi-cell wireless environment, TDMA uses so-called frequency planning, which means that neighboring cells do not operate in the same frequency band. CDMA uses frame expansion factor through direct spreading, in order to avoid frequency planning.

Error Correction Properties: In the case of CDMA, there is some correction of chip errors due to noise on the channel. The errors corrected are independent of the active user number

$A \leq n$. Here n is the user number upper limit. When A is n , then there is full occupancy, or in other words a maximum loading situation.

In the case of TDMA, error correction due to noise in the communication channel is not possible, irrespective of whether or not the channel is operating at less than full loading. There is then no error-correction gain from an increase in the upper user number limit n , as there is for the competing technology CDMA.

The Hamming distance between the code words is a key determinant of the error correction capability of a linear block-coding scheme. The greater this distance, the greater the error correction capability. This distance is at a minimum of 2 for TDMA with $n > 1$, and at a maximum of $\ell/2$ for CDMA $n \times \ell$ block codes where $\ell \geq n$ is a power of 2, that is 2, 4, 8, 16, 32.... Typically, direct sequence, orthogonal CDMA codes are square with $\ell = n$, and followed by *pseudo*-random CDMA codes.

Peak-to-Average Power Property: The peak power permitted to be transmitted in a communication cell is limited, either by law or other considerations. The peak-to-average transmitted power ratio is denoted PAP. For TDMA the PAP is unity under full loading and is n/A otherwise, see Figure 4. A PAP of unity is desirable so as to maximize signal-to-noise ratio (SNR), given peak power limits.

In CDMA, the transmitted signal is the summation of A synchronized polar, binary signals. It is termed an A -ary signal. In CDMA the peak-to-average power ratio is A . This ratio increases with the number of active users accommodated in the channel. Note that in theory, this increase limits the number of users n that can simultaneously use any CDMA wireless communication cell, but in practice, power control algorithms reduce the actual PAP of a CDMA system.

Computational Effort: Computational effort limitations in, for example, a mobile or low-cost receiver are more critical than in a hub station transmitter. Computational effort also affects receiver power consumption and battery life at the margins. CDMA systems require vector integer multiplication and sparse matrix integer multiplication operations. The corresponding calculations for TDMA are trivial.

Majority Logic Codes: There are nonlinear codes, termed here majority logic codes, which have application to CDMA coded signals according to embodiments of the present invention in communication systems. To multiplex n users, code word lengths of $\ell \geq 2^n - 1$ are proposed. With $n = 3$, $\ell \geq 7$; $n = 4$, $\ell \geq 15$; $n = 5$, $\ell \geq 63$; $n = 6$, $\ell \geq 127$; $n = 7$, $\ell \geq 255$, and so on. The frame expansion factor α becomes unrealistically high for applications as n increases, because of the associated system capacity loss. Such codes are not yet exploited significantly in the market place, and indeed recent research by others has concluded that their best future, in the absence of a major breakthrough, could be in niche markets with low user numbers. Their attractive property for applications, is that the PAP value is unity, as in TDMA, see Figure 4. Also there is some error correction at less than full occupancy, although there can be error introduction for some levels of occupancy. There is simplicity of implementation since the non-linearity in the code consists of merely a *sign* operation on the output of a linear code. This can be viewed as counting ‘yes’ or ‘no’ votes, so the majority logic aspect is simply a majority vote counting.

Most work on majority logic codes has been done for the case of square and single-stage codes for odd numbers of active users up to 7 users, mapping binary message data to binary transmission data. When not all users are active there can be error correction. However, there can also be deterministic errors. That is, errors are created, even for transmission over a noise free

channel. This occurs for the case of 5 users and 2 inactive users. In the case of $A \leq n = 7$, certain codes have been proposed and studied which use code words of length greater than 7, and although error correction improves, there are still errors in the noise free case for some occupancy levels. Such errors preclude wide acceptance of single stage Majority Logic Coding as it stands.

There is a notion of a *two-stage Majority Logic Coding*. Simulations have been carried out, for example, on 3×3 codes to achieve a scheme for odd order active users up to 9 users. Standard majority logic coding is first applied to three sets of 3 users. Then the 9 output signals are reordered ready for a further application to 3 lots of 3 signal sets. The decoding is a reverse double majority logic coding process. Figure 5 depicts the situation.

Majority logic coding, is a nonlinear coding, for which there is no complete theory. The generation of majority logic codes is a problem which appears non-polynomial (NP) hard: The computational effort to exhaust all possibilities grows at least at the rate of the order of $2^{(n \times n)}$, which is 2 to the power n all squared. Notwithstanding this, applicants have discovered that Go-CDMA codes are a class of codes which realize the potential of a nonlinear logic coding approach.

A Hadamard matrix is a matrix $H(n)$ with elements in the set $\{-1, +1\}$ such that $H(n)' \times H(n) = n \times I(n)$. A square $n \times n$ real matrix $H(n)$ for n exist if n is a factor of 4. This includes the *Hadamard* square $n \times n$ real matrices $H(n)$ for n a power of 2. That is, n belongs to a set denoted \mathbf{N} consisting of elements $\{2^m\}$ where m is the set of positive integers $\mathbf{Z}^+ = \{1, 2, 3, 4, \dots\}$. Thus

$$\mathbf{N} = \{n = 2^m \mid m \in \mathbf{Z}^+\}, \text{ or } \mathbf{N} = \{2, 4, 8, 16, 32, \dots\}. \quad (1.0)$$

These *Hadamard* matrices $H(n)$ for n a power of 2 can be simply constructed from a recursion involving the *Kronecker* product of *polar, binary-form* matrices with elements in the set $\{-1, +1\}$ as follows:

$$H(2^{(m+1)}) = H(2) \otimes H(2^m), \quad H(2) = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}. \quad (1.1)$$

- 5 This *Kronecker* product operation \otimes , replaces a $\{+1\}$ in $H(2^m)$ by $H(2)$ and a $\{-1\}$ by $-H(2)$. For example,

$$\begin{aligned} H(4) &= H(2^{1+1}) = H(2) \otimes H(2^1) \\ &= \begin{bmatrix} H(2) & H(2) \\ H(2) & -H(2) \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{bmatrix}. \end{aligned} \quad (1.2)$$

More generally, recall that with $X = (x_{ij})$ a matrix of scalar elements x_{ij} at the intersection of the i^{th} row and j^{th} column, then $X \otimes Y$ is a matrix built from X with x_{ij} replaced by the matrix $x_{ij}Y$.

The *Hadamard* matrix $H(n = 2^m) = (h_{ij})$ is then an $(n \times n)$ matrix with elements h_{ij} taking values in the set $\{-1, +1\}$; that is, $H(n)$ is a matrix in *polar, binary-form*. Let h_i^r denote the i^{th} row of H and h_i^c denote the i^{th} column of H , then $H = [h_1^c, h_2^c, \dots, h_n^c]$, $H' = [h_1^r, h_2^r, \dots, h_n^r]$. Here

15 the prime denotes the matrix transpose. That is, $H' = (h_{ij})$, or equivalently, to form H' , elements h_{ij} in H are replaced by h_{ji} for all i, j .

The *unipolar, binary-form* of a polar matrix, with elements in the set $\{0, 1\}$, is obtained by first adding unity to all of the polar matrix elements, and then dividing the result by 2. The polar form of the matrix is recovered by multiplication by 2 and then subtracting unity from all elements.

5 The matrix rows of a unipolar binary matrix $H(n)$ can be viewed as binary numbers each with n binary digits (reading left to right say). These have decimal equivalents. Likewise for the columns, or more simply the rows of the matrix transpose. Thus $H(n)$ can be represented by a set, or sequence, of n decimal numbers.

10 The *Hadamard* matrix H is *orthogonal*, in that $H' H$ is n times the identity $n \times n$ matrix $I(n)$. That is, $H' H$ has n diagonal elements that are all n , and there are zeros elsewhere.

Coding using Hadamard Matrices is now described. This is the basis of first-stage orthogonal-CDMA coding. Consider a data column n -vector d data given discrete time. For simplicity of presentation, assume that each element d_i of d is generated by a particular user, termed the i^{th} user; there being a set of n possible users. The data from an active user is $\{-1\}$ or $\{+1\}$, and from an inactive user is set as $\{0\}$. In general then, d is a ternary data vector.

In Hadamard matrix coding, the data transmitted is actually an n -vector which results from standard vector-matrix multiplication as follows:

$$s = Hd. \quad (1.3)$$

20 Thus the i^{th} element of s is $s_i = \text{sum of scalar products } (h_{ij} d_j)$ over all j . It is clear that with A active users, s_i is a signed integer in the range $\{-A, +A\}$ and results from integer multiplication.

The recovery of d from s is possible by the inverse operation,

$$d = \frac{H' s}{n} = \text{sign}\left(\frac{H' s}{n}\right), \quad (1.4)$$

where the *sign* of a vector is the sign of each element, as follows:

$$\text{sign}(x) = \begin{cases} +1 & \text{if } x > 0 \\ 0 & \text{if } x = 0. \\ -1 & \text{if } x < 0 \end{cases} \quad (1.5)$$

Alternatively, non-conventional definitions of a *sign* function can be given as follows:

$$\text{sgn}(x) = \begin{cases} +1 & \text{if } x \geq 0 \\ -1 & \text{if } x < 0 \end{cases}, \text{ or } \text{sgn}(x) = \begin{cases} +1 & \text{if } x > 0 \\ -1 & \text{if } x \leq 0 \end{cases}. \quad (1.6)$$

Note that (1.6) are the preferred definitions of the *sign* operation in the present embodiments of Go-CDMA. While the definition in (1.6) is not new, its application, instead of (1.5), in multi stage majority logic coding, is new and may be implemented according to embodiments of the present invention. The nonlinear operation in (1.5) would result in a ternary output but the output from either operation in (1.6) would be always binary.

The recovery of d follows since $H' s = H' H d = nI(n)d = nd$, by virtue of the orthogonality of H . Also, the *sign* operation on a vector with elements in the set $\{-1, +1\}$ introduces no change. In a multi-user system, the i^{th} receiver decodes only the element d_i using only one code word assigned to the i^{th} user, namely the i^{th} row vector of H' . That is, the i^{th} column vector of H , denoted h_i^c . Thus, the decoding for the i^{th} user extracts the scalar d_i as

$$d_i = \frac{h_i^{c'} s}{n}. \quad (1.7)$$

In the case of error-free transmission, then d_i belongs to the ternary set $\{-1, 0, +1\}$. However, in the case of transmission errors, then for an active user-receiver pair, $h_i^c s/n$ may not be in the set $\{-1, +1\}$, where $\text{sign}(h_i^c s/n)$ is in this set. Thus an estimate of d_i , denoted \hat{d}_i is

$$\hat{d}_i = \text{sign}\left(\frac{h_i^c s}{n}\right) \quad (1.8)$$

5 *Hadamard* codes are a special case of the class of orthogonal codes, since in the above coding, H can be replaced by any orthogonal matrix A with the property AA' proportional to I . The advantage of *Hadamard* codes is that their elements belong to the set $\{-1, +1\}$, so multiplications are trivial, and their Hamming distance (to be defined below) is at a maximum for such orthogonal codes.

10 If orthogonal codes with elements, of *unipolar binary form*, in the set $\{0, 1\}$ are used, and the data from active users are also in *unipolar binary form*, then the above mentioned matrix coding and decoding process will be a logical exclusive-OR operation instead of the standard vector-matrix multiplication. The standard vector-matrix multiplication approach is usually used, where both the orthogonal codes and active-user data are in the polar binary form.

15 Hamming Distance and Error Correction: The *code words* of a linear block code, such as H , are the rows, or columns, of H . The “distance” between two rows, or columns, of a binary code is the number of elements in one row that differ from the corresponding element in the other. The Hamming distance is the minimum distance of all the possible distances between rows or between columns. For a *Hadamard* code $H(n)$, with $n \in \mathbf{N}$, the Hamming distance is

$d_{\min} = n/2$. Standard theory on Hamming distances shows that $(d_{\min}/2) - 1$ errors are corrected in decoding.

Pseudo-Random Coding: *Pseudo*-random coding follows the same pattern as orthogonal coding, save that the elements of the orthogonal code matrix, denoted P , are generated by *pseudo*-random sequences. Typically, P is a rectangular code $(\alpha n \times n)$ -matrix with an integer frame expansion factor α . It is block diagonal with identical column blocks, being a *pseudo*-random code word vector of length α . For such coding

$$s = Pd, \quad \hat{d} = \text{sign}(P's) = \text{sign}(P'Pd) = d. \quad (1.9)$$

In the case of added channel noise, denoted *noise*, then

$$\hat{d} = \text{sign}(P'(s + \text{noise})) = \text{sign}(d + P'\text{noise}), \quad (1.10)$$

One virtue of this coding is that the term $P'(\text{noise})$ approaches an average of noise terms as α becomes large, which then diminishes the relative effect of the noise, with respect to the signal. This is the case, irrespective of the noise characteristics, so becomes important in multi-cell applications. In this situation, the noise includes interference from neighboring cells. These use different P matrix codes, denoted P_i , so that $P'P_i$ is relatively small compared to $P'P$.

TDMA as coding: In TDMA systems, the H is replaced by the identity, so $s = I(n)d$, and $d = I'(n)s = I'(n)I(n)d = d$. The i^{th} user “decodes” trivially the received signal s to achieve $d_i = s_i$, the i^{th} element of s .

In TDMA, there is frequency division multiplexing of cell signals in order to reduce interference from neighboring cells. The available spectrum is divided into β frequency bands, typically $\beta = 4$, or so, and neighboring cells are allocated to different bands. There is a process

termed *frequency planning*, to allocate the frequency bands. The avoidance of this process by using *pseudo-random-CDMA* coding, in CDMA systems, is considered a major advantage of CDMA. In order for the receiver to properly decode signals received from transmitters that implement pseudo-random-CDMA coding, the receiver must be aware of the pseudo-random implementation. This information in the form of the permutation matrix may be transmitted from the transmitter to the receiver via a pilot CDMA channel or via a data channel. Alternatively, all pertinent permutation matrices can be pre-stored in a memory or storage module in the communication device. In this implementation the information sent from the transmitter to the receiver via a pilot CDMA channel or via a data channel is only for locating, within the memory or storage module of the communications device, the correct permutation matrix to be used.

Majority Logic Coding: In conventional majority logic coding using square block (nonlinear) codes, the *Hadamard* matrices $H(4)$, or $H(8)$, are reduced in size by 1, by a deletion of the first row and first column, consisting of only elements $\{+1\}$. This achieves square majority logic coding matrices, denoted $M(3)$ or $M(7)$. Thus, the deletion serves to permit a clear majority voting, possible only for the case of odd numbers, and not for the case of even numbers, as in the set \mathbf{N} . Then Majority Logic Coding using $M(7)$ (or $M(3)$) encodes a 7-vector (or 3-vector) d as

$$s = \text{sign}(Md), \quad (1.11)$$

using the *sign* operation of (1.5). The result is a ternary vector, rather than binary vector. The coding situation is depicted in Figure 6. The decoding is done by the operation

$\hat{d} = \text{sign}(M'(s + \text{noise}))$. The term $(s + \text{noise})$ is the received signal vector with additive noise.

The i^{th} receiver with the code word m_i^c , being the i^{th} row of M decodes as

$$\hat{d}_i = \text{sign}(m_i^{c'}(s + \text{noise})). \quad (1.12)$$

The decoding situation is depicted in Figure 7. The PAP is unity.

5 There is a variation to Majority Logic Coding, where certain rectangular codes $(\ell \times n)$ -
matrices $M(\ell, n)$ are used for $\ell > n$. For example, with $n = 7$, the codes
 $M(n \times 2n) = [M(n) \mid +M(n)]$ and $M(n \times 2n) = [M(n) \mid -M(n)]$ have been tested to achieve
improved error correction. Also, the majority logic codes of length $\ell \geq 2^n - 1$, noted already, are
simply n arbitrary column selections from a Hadamard matrix of size 2^n , or with one row
deleted.

II. Go-CDMA Technology Overview

Go-CDMA technology is suited for implementation in CDMA communication systems. Go-CDMA coding is nonlinear, multi-stage coding for application in a hub station transmitter. Go-CDMA decoding in the receiver involves the inverse process. The Go-CDMA multi-stage codes are constructed from the interconnection of proposed Go-CDMA nonlinear, building-block codes. These in turn are constructed from Go-CDMA matrices, satisfying a proposed Go-CDMA matrix defining property. These constraints ensure attractive error correction properties for the system under partial loading conditions. In this subsection, a qualitative overview is presented and in the next subsection, more precise mathematical descriptions of the technology according to embodiments of the present invention are given.

Building-block Go-CDMA codes: The Go-CDMA building-block codes are rectangular nonlinear codes, constructed from rectangular Go-CDMA matrices satisfying Go-CDMA matrix defining properties. The matrices are extracted from the rows and columns of orthogonal Hadamard matrices. The coding requires linear integer arithmetic operations on the data using these matrices and certain *sign*, *sgn*, *Sgn* or *integer rounding* nonlinear operations. The *sign* operation is as defined in (1.5) while the *sgn* and *Sgn* operations are defined in (1.6) and (1.18), respectively.

Consider first the case of a simple one-block $\ell \times n$ -Go-CDMA code. This is for $A \leq n$ active users transmitting binary data in frames of $\ell = \alpha n \geq n$ chips, giving a frame expansion factor of $\alpha = (\ell / n) \geq 1$. The coding involves linear block coding of polar ternary vectors, which include the message data, followed by a *quantization* operation, such as a *sgn* operation, or when used in a multistage coding a variation of this operation termed here a *Sgn* operation. This yields

multi-level data for transmission. After the transmission of this data, the first stage of decoding involves linear block decoding of the multi-level data, followed by a *sgn* operation to recover the binary message data estimates.

The incoming vector information, without loss of generality, has elements in the ternary set $\{-1, 0, +1\}$. The i^{th} user data is in the i^{th} element of the message vector. If the user is active, then the data is polar-binary data with elements in the set $\{-1, +1\}$, and if the user is inactive, then the data is simply the set element $\{0\}$. If there is a pause in transmission, a sentinel code is usually transmitted with elements in the set $\{-1, +1\}$. The alternative of transmitting $\{0\}$, can cause noise in reception during a pause.

For the multi-cell communication system situation, by using different codes in different cells, there is some suppression of neighbor cell interference.

Go-CDMA coding, using one building-block code, is free of so-called deterministic errors, irrespective of the number of active user numbers. It corrects errors, diminishing approximately linearly from $\ell/2$ for ℓ even, or $(\ell/2)-1$ for ℓ odd, in the case $A=1$. The decrease is to a positive integer in the full loading case $A=n$. The positive integer is denoted $e(\ell)$, which also increases as ℓ increases, although not necessarily monotonically, being in the range $[2, 7]$ in building block codes with $n \leq 4$, $\ell \leq 16$. Examples of building-block code sizes with these preferred error correction properties are 8×3 , 16×5 , 12×5 , 20×7 and 24×7 .

Go-CDMA building-block codes, for specified frame expansion factors, achieve error correction, which increases linearly with the maximum user number n , and is computationally simpler per user as n increases. The upper limit on n is about $n=8$, or so, since beyond this,

the frame expansion factor α necessary to permit satisfaction of the Go-CDMA matrix defining property, becomes too large for application to multi-stage coding.

An increase in higher user numbers in single-stage Go-CDMA coding can be achieved by paralleling Go-CDMA building blocks, preserving error correction within the blocks, but with no possibility for cross block error correction, unless followed by later stages of Go-CDMA coding.

Two-stage Go-CDMA coding: Two-stage coding is a concatenation of two compatible, single-stage, parallel, block-coding stages with suitable interconnections. Preferably, the interconnections are designed to achieve a scaling up of the error correction capability of the building-block codes, or at least as best a scaling up as can be achieved. Such Go-CDMA codes, in the case of identical size building blocks, can be used conveniently for user numbers, being the square of the numbers handled by a single building-block code.

Consider a pair of $a \times b$ and $c \times a$ Go-CDMA single stage codes where each single-stage code is constructed by parallel connection of $a_1 \times b_1$, and $c_1 \times a_1$ building-block codes. Then in a two-stage interconnection to achieve a $c \times b$ code, the goal is to achieve for $1 \rightarrow b$ users, $c/2 \rightarrow e(c)$ errors corrected. Thus, as the user numbers are squared in a two-stage coding, the goal is that the errors corrected in a frame are also squared.

As an example, consider a 16×5 -Go-CDMA single-stage building-block code, for implementing in multi-stage coding. This block has a 5-vector input and 16-vector output. Interconnections of such building blocks may achieve two-stage 256×25 -Go-CDMA codes. The interconnection arrangements for this is that there are 5 parallel building blocks driven by a total of 25 inputs in a first stage coding. The five sets of outputs of this first stage, totaling 80 outputs, become inputs

to a second set of 16 parallel building blocks giving now 256 outputs. This is depicted in Figure 8.

The decoding process reverses the flow of information, so that the outputs, after transmission, become inputs for the decoding blocks. These pass through two stages of single stage decoding to recover the original message, in the absence of transmission channel noise. In reception, each receiver only decodes its assigned information. There are two rules for applying two-stage Go-CDMA coding.

Rule 1: No two outputs from a first-stage building block should become inputs to a single second-stage building block. As long as this rule is obeyed, the connections can be either systematic or *pseudo* randomized. This rule ensures that there is optimum error correction in the multi-stage coding so that there is a scaling up the error correction capability of a single building-block code.

Rule 2: Within the one communication cell, either the code selection from the set of building-block nonlinear codes of the same size is by a *pseudo*-randomized selection from a large class. One realization of this is to ensure that the interconnections between the different stages, using standard building blocks, are *pseudo* randomized. As long as this rule is obeyed, each cell using multi-stage Go-CDMA coding will have a “randomly” different code, so that the inter-cellular interference in the decoding process is reduced by the frame size expansion factor.

The above rules are not compulsory but may be followed to obtain optimal results from a two or more stage Go-CDMA coding scheme. If non-optimal results would suffice, then any invertible permutation operation can be used in a two or more stage Go-CDMA coding scheme.

5 Multi-stage Go-CDMA coding: Three-stage Go-CDMA coding is a concatenation of three compatible, single-stage codings with suitable interconnections. Each stage is constructed by paralleling suitably sized building block codes. Preferably, the interconnections are designed to achieve a true scaling up of the error correction capability of the building-block codes, and thereby achieve Go-CDMA codes for the cube of the user numbers that can be handled by a
10 single building block code, at least in the cases when all building blocks are identical. Consider $a \times b$, $c \times a$, $f \times c$ Go-CDMA single stage codes. Then in a three-stage interconnection to achieve an $f \times b$ code, the goal is to achieve for $1 \rightarrow b$ users, $f/2 \rightarrow e(f)$ errors corrected. This is illustrated in Figure 9.

15 Four-stage Go-CDMA coding is an extension of the two-stage and three-stage Go-CDMA coding approach. This can effectively raise the building block size by a power of 4, scaling up the error correction properties accordingly. Four-stage Go-CDMA coding can be implemented as two nested two-stage codes. This is illustrated in Figure 10.

20 By working with rectangular block codes resulting in different extended frame sizes, the users may be permitted different classes or quality of service. Some data could be more important than others for accurate transmission, so this user data can be assigned extended length codes.

In multi-cell wireless applications, selection of the frame expansion factor of typically, $\alpha = 64$ or higher, would be needed to give significantly “better” neighbor cell interference suppression than frequency planning.

Alternative Go-CDMA realizations: In alternate embodiments of the invention, larger building-block Go-CDMA codes are readily designed, which satisfy mildly relaxed Go-CDMA matrix defining properties. They have reduced error correction capability, and also introduce deterministic errors for some partial loadings, unless something is done to prevent this. If such codes are used, *pseudo*-data, perhaps just $\{+1\}$, may be generated and substituted for a possible small subset of inactive users prior to coding and transmission. This subset number is chosen to maximize the error correction. Such *pseudo* data can be used also to ensure an odd number of active users so that binary signals are transmitted, as in traditional majority logic coding, or to eliminate deterministic errors which arise for certain user numbers and using certain “less preferred” codes. Thus we may consider the case of n users, including both A active users and certain inactive users becoming *pseudo*-active, and the remaining users inactive. One example of a code requiring such *pseudo* active-user data to eliminate deterministic errors, is a 512×20 code.

With a full complement of *pseudo*-active users, as required if calculating with using exclusive-OR operations, there is an effective conversion of the raw data set of ternary elements in the set $\{-1, 0, +1\}$, with zero indicating an inactive use, to the polar binary set $\{-1, +1\}$. Figure 11 depicts the assembly of user data, including active user data, *pseudo*-active user data and inactive user data into an n -vector d .

When user number limits n , are prescribed *apriori* to be outside the sets of the integer powers of the integers ...3,4,5,6,7,8..., there are suitable realizations, but perhaps less preferred

because they appear unnecessarily complicated. One performs a multi-stage Go-CDMA coding using different code sizes in each stage, another includes inactive users to pad the numbers up to some n belonging to one of the preferred.

Error correction: For a building-block $n \times \ell$ Go-CDMA code, the underlying minimum distance (Hamming distance) between the Go-CDMA matrix columns is typically in the range $[2, (\ell+1)/2]$ and for the rows, 2, at least for a $n > 3$. Other measures are the singular values, or eigenvalues, of the Go-CDMA code matrix. These measures are invariant of the ordering of the rows and columns. Working with Go-CDMA matrices that achieve a maximum Hamming column distance close to $(\ell+1)/2$, and the minimum ratio of maximum-to-minimum singular value (close to unity) gives good error correction properties. As an example of error correction, for a particular 16×4 Go-CDMA code, there is a correction of 2 bit errors at full loading, increasing to 7 bit errors with $A=1$ active users. In multi-stage coding, the goal is to effectively scale up these error correction figures.

When there is less than full channel occupancy, the Go-CDMA coding is such that there is correction of some of the errors introduced by channel noise. Larger lengths ℓ of code words, by factors of 2,4,8,... can always be used to achieve greater channel error correction, and to reduce the relative strength of neighbor cell interference, but at a loss in capacity. Without neighboring cells, there is perhaps some incentive to have frame expansion factors of greater than unity. However, with neighboring cell interference, and in the absence of frequency planning, there is an imperative to have expansion factors of at least 40 to compete with TDMA which uses frequency planning.

In practical applications, as the allowable user number n increases in a communication cell, there is less chance of full or high levels of occupancy, and more error correction. Code

allocation across cells in a cluster of cells also decreases high occupancy levels and increases error correction. Go-CDMA has the potential to cope with larger user numbers per cell than CDMA.

Computational Effort: Single stage Go-CDMA coding and decoding have simplicity of implementation in that computational effort for single-stage, building block codes is of the same order as that for current linear block coding schemes of the same size, and upgrading is primarily a software or firmware change to change the linear codes and additional mathematical operations. For example, *quantization operations*, typically *sgn* operations, and *integer rounding* operations may be included to implement the Go-CDMA coding. Multi-stage Go-CDMA coding is likewise not really any more expensive than CDMA coding, although its decoding requires up to an order of magnitude computational effort increase for the user.

Majority Logic Coding goal: Go-CDMA achieves the goal of Majority Logic Codes by departing sufficiently from conventional Majority Logic Coding approaches to make it useful. The applicants have, for example, developed the “complete” set of Majority Logic Codes, which introduce no deterministic errors, at least for the case of realistic frame expansion factors. This set is relatively sparse. The applicants have also generated a set of Go-CDMA codes which introduce no deterministic errors with all users active, but otherwise may introduce errors. These sets of codes include the earlier studied incomplete set of codes for $n = 3$, $n = 7$, and the codes with word length $\ell \geq 2^n - 1$. To compete with present day CDMA and TDMA, codes of order 256 are desirable. Such codes are available as Go-CDMA codes. To compete with TDMA in terms of inter-cell interference, Go-CDMA codes having frame expansion factors $\alpha \geq 64$ are preferable.

III. Go-CDMA Matrices, Coding, Decoding and Preferred Embodiments

The Go-CDMA matrix property is defined below in various ways including through the use of equations and accompanying description. These definitions are used to specify what matrices are considered Go-CDMA matrices herein. Codes that comprise of rows and/or columns of the Go-CDMA matrices are considered Go-CDMA codes. Consider an $a \times b$ -matrix $M(a, b)$, or denoted M , with $a \geq b$, and having real finite elements. Consider also data b -vectors d , with elements d_i in the ternary set $\{-1, 0, +1\}$. Then this is denoted a Go-CDMA matrix if,

$$d_i = \text{sgn}(m_i^c ' Q_k(Md)) \text{ for all } i \in \{i \mid d_i \in \{-1, +1\}\}, \quad (1.13)$$

or, in the case that $k = 1$,

$$d_i = \text{sgn}(m_i^c ' \text{sgn}(Md)) \text{ for all } i \in \{i \mid d_i \in \{-1, +1\}\}, \quad (1.14)$$

where m_i^c is the i^{th} column of M . In the above defining equation, there is a quantization operation $Q_k(\cdot)$. This is a standard operation that converts an input integer in the set $[-a, +a]$ to an output belonging to the set $[-k, -1] \& [+1, +k]$ for some integer $k \leq a$. The conversion formula allows arbitrary mappings, but these must be specified. For example, the integer space can be “halved” in that input integers in the range $[-2k, 2k]$ could map into the above set by first halving each integer, and using an integer rounding operation. Thus in obvious notation, apply $\text{round}(Md/2)$. Here $\text{round}(r) = \text{sgn}(r)$, for $-1 \leq r \leq 1$. Multiple “halvings” can also be used easily. Also, integers in the set $[-qk, qk]$ where q is an integer, can be mapped into the desired set by first dividing each integer

by q followed by a rounding operation. Note that $sgn(\cdot) = Q_1(\cdot)$. Also, note that if $k = a$, then $Q_a(Md) = Md$, and the operation is trivial and linear for $Md \neq 0$.

Any code matrices which can satisfy the operations defined in (1.13) and/or (1.14) is considered
5 Go-CDMA codes herein.

One-stage Go-CDMA coding: One stage of Go-CDMA coding using a Go-CDMA code matrix M is,

$$s = Q_k(Md), \quad (1.15)$$

where s is the transmitted signal a -vector, in general ternary, and the sgn and *quantization* operations are as in the Go-CDMA matrix defining property. The signal s is $(2k)$ -ary. With $k = 1$, $Q_1(Md) = sgn(Md)$, and s is ternary.

One-stage Go-CDMA decoding: The corresponding Go-CDMA decoding stage is

$$\hat{d} = sgn(M'r), \quad (1.16)$$

where r is the received signal after transmission and is thus $s + noise$, where the notation *noise* stands for additive transmission channel noise.

Consider Go-CDMA matrices $M = M_1(a, b)$, $M_2 = M_2(c, a)$. Then we can define a two-stage
20 Go-CDMA coding as the application of a single stage Go-CDMA coding, followed by a second single stage Go-CDMA coding.

Two-stage Go-CDMA coding: For Go-CDMA matrices $M_1 = M_1(a, b)$, satisfying (1.14) and $M_2 = M_2(c, a)$, satisfying (1.13), and a permutation matrix P , two stage of Go-CDMA coding is,

$$s = Q_k(M_2 P Sgn(M_1 d)), \quad (1.17)$$

- 5 where s is the transmitted signal c -vector, which is ternary or $(2k+1)$ -ary. Notice that the quantizer is only applicable in the second stage. Also, with $k=1$, $Q_1(Md) = sgn(Md)$, and s is ternary. Notice also that instead of a sgn operation on the vector $M_1 d$, a modification is used, denoted Sgn defined as follows:

$$Sgn(x) = \begin{cases} sgn(x) & \text{if } x \neq 0 \\ 0 & \text{if } x = 0 \end{cases} \quad (1.18)$$

where x is a vector.

Two-stage Go-CDMA decoding: The corresponding Go-CDMA decoding stage is,

$$\hat{d} = sgn(M_1' P' sgn(M_2' r)). \quad (1.19)$$

- 15 Three-stage Go-CDMA coding: Three-stage Go-CDMA coding using Go-CDMA matrices M_1, M_2 , satisfying (1.14), and M_3 satisfying (1.14) or (1.13), and permutations P_1, P_2 of appropriate dimensions is,

$$s = Q_k(M_3 P_2 Sgn(M_2 P_1 Sgn(M_1 d))), \quad (1.20)$$

- where s is the transmitted signal c -vector, and k is some integer satisfying $k \geq 1$. The quantizer
20 is only applicable in the last stage.

Three-stage Go-CDMA decoding: The corresponding Go-CDMA decoding stage is,

$$\hat{d} = \text{sgn}(M_1' P_1' \text{sgn}(M_2' P_2' \text{sgn}(M_3' r))), \quad (1.21)$$

Multi-stage Go-CDMA coding and decoding: For four stage and a higher number of stages, the same pattern for generalization follows that now established.

The Go-CDMA matrix property ensures that there are no deterministic errors for any active users $A \leq n$ in the nonlinear coding and decoding processes. There may be errors for inactive users when $d_i \in \{0\}$, but there is no destination for such errors.

Go-CDMA matrices $M(a,b)$ as signed permutation matrices: One class of Go-CDMA matrices according to an embodiment of the present invention is the class of signed permutation matrices. A signed permutation matrix has all elements zero, save that in each row and column there is one, and only one, element belonging to the set $\{-1, +1\}$.

Go-CDMA matrices $M(a,b)$ from Hadamard matrix columns: In Table 1, there are listed Go-CDMA code matrices $M(a,b)$ consisting of certain b columns of the Hadamard matrix $H(n)$ and extracting the a rows. Such Go-CDMA matrices partially inherit the Hamming distance properties, depending on how many rows and columns are preserved, and the particular selections. Note also the errors corrected at full loading e in Table 1.

Go-CDMA matrices $M(a,b)$ as block-diagonal Hadamard sub-matrices: If $M_i(a_i, b_i)$ are Go-CDMA matrices for all $i \in \{1, 2, \dots, q\}$, then the block diagonal matrix

$$M(a,b) = \text{block diag}\{M_1(a_1, b_1), M_2(a_2, b_2), \dots, M_q(a_q, b_q)\}, \quad (1.22)$$

5 is also a Go-CDMA matrix, where $a_1 + a_2 + \dots + a_q = a$, $b_1 + b_2 + \dots + b_q = b$. Of course, the simplest case is when the a_i are all equal to a_1 , and the b_i equal to b_1 . A special case then is if $q = a_1$, so that $M(a,b) = M(a_1^2, ab)$.

Go-CDMA augmented matrices $M(a,b)$: If $M_i(a_i, b)$ are Go-CDMA matrices for all $i \in \{1, 2, \dots, q\}$, then the augmented matrix

$$M(a,b) = [M_1(a_1, b)', M_2(a_2, b)', \dots, M_q(a_q, b)']', \quad (1.23)$$

is also a Go-CDMA matrix. Of course, the simplest case is when the $M_i(a_i, b)$ are all equal.

Class of Go-CDMA code matrices: There is a class of codes $M_{class}(a,b)$ satisfying the Go-CDMA matrix property above, constructed from one such code $M(a,b)$, by all possible re-
 15 ordering of rows and columns of $M(a,b)$ and all possible column and row sign changes. That is

$$M_{class}(a,b) = \left\{ \left(P(a) M(a,b) \right) P(b) \mid \begin{array}{l} P(m) \text{ is the set of all signed} \\ m \times m \text{ permutation matrices} \end{array} \right\}. \quad (1.24)$$

Defining the code words of $M(a,b)$ as the columns of $M(a,b)$, together with the columns of $-M(a,b)$. The complete set $M_{class}(a,b)$ can be constructed by all possible ordered k code
 20 word selections from this set of code words, save that if one code word is selected, then its

negative is excluded. There are $2^b b!$ distinct member codes: For $b=4$, this number is $16 \times 24 = 384$, and for $b=8$, this number is $(256) \times (40,320) = 10,321,920$.

n	Rows a	Columns b	Errors Corrected e	Columns preserved in $H(n)$
32	1, 4, 5, 8, 9, 12, 13, 16, 17, 18, 20, 23, 24, 26, 27, 28, 30, 32	≤ 5	1	3, 6, 9, 22, 27
16	1 to 16	≤ 5	2	2, 7, 14, 25, 27
12	1 to 12	≤ 6	0	5, 6, 7, 9, 10, 12
8	1 to 8	≤ 3	1	2, 3, 8

Table 1. Go-CDMA matrices $M(a,b)$ from the rows and columns of $H(n)$.

Go-CDMA coding-decoding from one basic code matrix M according to a preferred embodiment of the invention: Consider a Go-CDMA code matrix $M(a,b) \in M_{class}(a,b)$ constructed from a Hadamard matrix H . Then construct a Go-CDMA building-block code from this matrix. In turn construct a two-stage, three-stage, or four-stage Go-CDMA coding block using the one building block. Form block diagonal matrices from $M(a,b)$ as $M(a^2, ab), M(ab, b^2)$ ready for two-stage coding and decoding. Or $M(a^3, a^2b), M(a^2b, ab), M(ab, b^3)$, for three-stage coding and decoding. Or from $M(a^4, a^3b), M(a^3b, a^2b^2), M(a^2b^2, ab^3), M(ab^3, b^4)$ ready for four-stage coding and decoding.

The selection of $M(a,b)$ from $M_{class}(a,b)$ for each diagonal block are the same for all users in a cell, but are *pseudo*-randomized across cells is to ensure interference reduction from neighboring communication cells transmitting in the same frequency spectrum.

Between each stage of coding (and decoding) and the next stage, there may be a “mezzanine” stage, which is a permutation operation using a P matrix on the outputs of the each stage to re-order the outputs for inputting to the next stage. Thus in a two stage scheme, for the coding process the columns of an $a \times b$ matrix are formed by the vector outputs of each first stage coding, and the rows of the matrix extracted and concatenated to form one row vector. Of course in the decoding process, the inverse permutation matrix is used. That is, the vector inputs become rows of a matrix from which columns are extracted in order to form the vector outputs.

Best stage number, user-number limits and frame expansion factor: The most widely applicable multi-stage Go-CDMA coding for larger user numbers, less than 64, is two-stage coding. For larger user numbers, three-stage coding may go to 512 users, and four stages may be implemented to go higher. One-stage coding is limited in its error correction capability, but can be used for any user number limits. For n -user code-division multiplexing applied to mobile communications, there is an elegance in the user number limits n being either 64, 512, 625, 1296, or other squares, cubes or fourth powers of integers up to 8. Compare this to the current conventional CDMA usage of 64, 128 and 256.

The Go-CDMA technology holds out improved performance as n increases. Therefore, according to one embodiment of the invention, n should be maximized subject to other system design constraints.

The best frame expansion factor α depends on the expected level of neighbor cell interference. With zero such interference, as in frequency planning, it may be best to take a small factor given by minimum-length code words. Otherwise, to compete with TDMA and frequency planning, code word length may be increased to achieve desired factors of typically above 40.

5 Examples are selected to achieve minimum α of less than 128, but higher factors are readily obtained by working with a higher α value in the selected building block matrices $M(a, b)$, for a specified b . Also, one can work with higher order Hadamard matrices, or by building composite Go-CDMA matrices.

Pre-coding and post-decoding of signals: To improve transmission bit error rates in high channel noise, as in all communication systems for such channels, there are advantages in applying conventional pre-coding and post-decoding of signals.

Alternative embodiments for carrying out the invention: A generalization of the Go-CDMA matrix defining property (1.13) is to allow W -ary data in the data vector d , for $W > 3$, and to use an appropriate quantization operation in a generalized defining property. This may be less preferred in the current applications environment because the complexity may be unwarranted for applications currently envisaged. However, it does allow the Go-CDMA coding and decoding to generalize all its *sgn* operations to *quantization* operations to give further options.

Replacing the *sgn* or *Sgn* operations with *sign* operations in the Go-CDMA matrix defining property, except at the last decoding stage, is an alternative. In this case, care must be exercised in second or later stage coding, since zero-element input data no longer represent inactive users, so that such zero elements do not introduce errors.

Another embodiment which is a compromise is to relax the Go-CDMA code matrix requirement.

Consider an $a \times b$ matrix M , or $M(a,b)$, and data n -vectors d , with elements d_i in the binary

set $\{-1, +1\}$, then choose matrices that satisfy: $d = \text{sgn}(M' \text{sgn}(Md))$ for all d . Now there are no

deterministic errors when all users are active, but there can be errors otherwise. Thus, there is a

5 need to add *pseudo*-active users when necessary to achieve no deterministic errors. Useful

specific codes, derived from Hadamard matrices $H(a)$ are listed in Table 2, if such a

compromise is considered seriously:

<u>m</u>	<u>a</u>	<u>b</u>	Preserved columns in $H(2^m)$
4	16	9	1 to 5, 7, 9, 10, 13
4	16	10	1 to 5, 7, 9, 10, 13, 16
5	32	11	1 to 11
5	32	12	1 to 11, 17
6	64	13	1 to 12, 17
6	64	14	1 to 12, 17, 33
7	128	15	1 to 13, 17, 33
7	128	16	1 to 13, 17, 33, 65
8	256	17	1 to 14, 17, 33, 65
8	256	18	1 to 14, 17, 33, 65, 129
9	512	19	1 to 15, 17, 33, 65, 129
9	512	20	1 to 15, 17, 33, 65, 129, 257
10	1024	21	1 to 15, 17, 33, 65, 129, 257, 513
10	1024	22	1 to 17, 33, 65, 129, 257, 513

Table 2: Go-CDMA matrices $M(a,b)$ with relaxed constraints. Here $a = 2^m$.

10

The pattern above enables the ready extension of these results to higher b values.

Another option is to achieve additional frame expansion by conventional post-*pseudo*-random-CDMA coding along with pre-*pseudo*-random-CDMA decoding.

Other embodiments may use alternative sign function operations, or exclusive-OR operations in the Go-CDMA coding context to achieve effectively the same or similar properties.

Fig. 12 depicts a method of coding signals using Go-CDMA codes for transmission within a CDMA system. Referring to Fig 12, in step 400, the multiple data messages enter a coder. The data messages typically have been pre-processed or pre coded to include error correction or protection bits. The data messages may be received from a processor or other hardware within a CDMA device or a software routine. The coder may be implemented in hardware or software that is the same as or distinct from the hardware and/or software that produced the data messages as illustrated and described with reference to Figure 3A.

In step 410, psuedo active user data is optionally inserted as a data message into the data message stream for coding. This step may be performed at the illustrated point in the process or may be performed prior to step 400 by a processor or in steps 430 or 440. This step may be performed when desirable depending on the number of active users within a system, the total number of permitted users and the Go-CDMA codes chosen in order to prevent deterministic errors within the CDMA system.

Step 420 determines the coder implementation. Step 420 need not be an actual processing step within a Go-CDMA coder implementation. If the coder is implemented as a single stage coder, then step 430 begins. If the coder is implemented as a multi-stage coder then step 440 begins.

In step 430, the coder codes the received data messages, and any pseudo active data messages, based on Go-CDMA codes. The coding may be performed, for example, according to equation (1.15). The coded data is then made available for modulation in step 450. In step 440, the coder codes the received data messages, and any pseudo active data messages, based on Go-

CDMA codes and multi-stage techniques. The coding may optionally include inserting pseudo-randomized connections between stages. When two stage coding is implemented, equation (1.17), for example, may be used. When three stage coding is implemented, equation (1.20), for example, may be used. When more than three stage coding is implemented, the equation that is implemented may be extrapolated based on the pattern established among the equations for one, two and three stage coding.

In step 450, the data messages encoded based on the Go-CDMA codes is channel modulated. Then in step 460, the modulated signal is transmitted over a communications channel. In the context of a base station transmitter, the data messages that are fed into the coder may be multiplexed data message streams corresponding to many active users and the communications channel is generally air. In the context of a mobile unit transmitter, the data messages that are fed into the coder may be for a single data message stream or a few message streams. When the coder implements the optional pseudo-randomized connections, the transmitter may also transmit pseudo-randomized connection data regarding the pseudo-randomized connection scheme to decoders to facilitate decoding. The pseudo-randomized connection data may be data that characterizes the scheme or data that identifies it to a remote device that includes multiple stored pseudo-randomized connection schemes stored in memory.

Fig. 13 depicts a method of decoding signals using Go-CDMA codes received within a CDMA system. Referring to Fig 12, in step 500 a receiver receives a composite CDMA signal from a communication channel. In step 510, a demodulator demodulates the composite signal to recover multiplexed data messages and conveys the multiplexed data messages to the Go-CDMA code based decoder.

Step 520 determines the decoder implementation. Step 520 need not be an actual processing step within a Go-CDMA decoder implementation. If the decoder is implemented as a single stage decoder, then step 530 begins. If the decoder is implemented as a multi-stage decoder then step 540 begins.

5 In step 530, single stage decoding is used to decode one or more data streams based on Go-CDMA codes. The decoding may be implemented, for example, according to equation (1.16). The decoded data is then made available to a post decode or other unit for further processing in step 560.

10 In step 540, when psuedo-randomized connections are optionally used during the coding process, the decoder is configured to have the corresponding pseudo-randomized connections necessary to decode the signal. This pseudo-randomized connections within the multi-stage decoder may be configured based on pseudo-randomized configuration data received from the transmitter within the CDMA system and/or based on pseudo-randomized configuration data stored in memory.

15 In step 550, multi-stage decoding is used to decode one or more data streams based on Go-CDMA codes. When two stage decoding is implemented, equation (1.19), for example, may be used. When three stage decoding is implemented, equation (1.21), for example, may be used. When more than three stage decoding is implemented, the equation that is implemented may be extrapolated based on the pattern established among the equations for one, two and three stage
20 decoding. The decoded data is then made available to a post decode or other unit for further processing in step 560.

In step 560, the decoded data is received for further processing which may include error correction based on error correction codes within the decoded data messages.

Figure 14 depicts an illustrative method of generating Go-CDMA codes according to an embodiment of the present invention. Referring to Fig. 14, in step 600, n codes of length ℓ are selected. In step 610, it is determined whether the codes satisfy equation (1.13) or (1.14). If not, then step 600 is repeated. If so, then in step 620 the codes are stored as Go-CDMA codes. It may be convenient to qualify $n \times \ell$ blocks at the same time, or portions thereof. Notwithstanding the foregoing, larger or smaller blocks may be qualified at the same, particularly in the case of multi-stage coding and decoding, and the steps of Figure 14 are merely meant to be illustrative of a technique for generating codes. The processes for using Go-CDMA codes, in general, do not depend on the generation method.

While specific embodiments of the present invention have been disclosed, it will be understood by those having ordinary skill in the art that changes may be made to those embodiments without departing from the spirit and scope of the invention. It will be further understood that the mathematical and matrix operations using Go-CDMA codes and matrices may be implemented in hardware or software. In the latter case, software instructions and data may be embodied in a computer useable medium and stored in a memory of a communications device. The software instructions may include control logic which when executed by a processor or other hardware cause the communications device to encode and decode data messages based on the Go-CDMA codes as depicted in Figures 4 – 13 and described above, or to generate Go-CDMA codes as depicted in Fig. 14. When implemented in hardware or firmware, the mathematical and matrix operations using Go-CDMA codes and matrices may be provided by logic on one or more chips or may be burned into, for example, a EEPROM as program instructions and data.

CLAIMS:

What is claimed is:

1. A method for coding a code division multiple access signal based on Go-CDMA codes,
comprising:
 - 5 providing majority logic coding blocks, each block comprising a Go-CDMA matrix;
coding a data message based on the majority coding blocks; and
transmitting the coded data message over a communication channel.
- 10 2. The method according to claim 1, wherein the majority coding blocks comprise a single coding stage.
3. The method according to claim 1, wherein the majority coding blocks comprise at least two coding stages.
- 15 4. The method according to claim 1, wherein the majority coding blocks comprise at least three coding stages.
5. The method according to claim 1, wherein the majority coding blocks comprise more than three coding stages.
- 20 6. The method according to claim 1, further comprising coding a plurality of data messages based on the majority coding blocks.

7. The method according to claim 6, wherein at least one of the data messages is associated with an active user.

8. The method according to claim 6, wherein at least one of the data messages is associated with a pseudo active user.

9. The method according to claim 6, wherein at least one of the data messages is associated with an inactive user.

10. The method according to claim 3, further comprising a permutation stage between each of at least one adjacent pair of the coding stages.

11. The method according to claim 4, further comprising a permutation stage between each of at least one adjacent pair of the coding stages.

12. The method according to claim 5, further comprising at least one permutation stage between each of at least adjacent pair of the coding stages.

13. The method according to claim 1, wherein the data message comprises data message elements in polar binary format.

14. The method according to claim 1, wherein the data message comprises data message elements in ternary format, or higher order formats.

15. The method according to claim 6, wherein each of the plurality of data messages is based on data received from an intermittent data source.

5 16. The method according to claim 1, wherein the majority coding logic blocks are implemented as a look up table and wherein the coding is performed based on the look up table.

17. The method according to claim 2, wherein the majority coding logic blocks are implemented as a look up table and wherein the coding is performed based on the look up table.

10 18. The method according to claim 3, wherein the majority coding logic blocks are implemented as a look up table and wherein the coding is performed based on the look up table.

15 19. The method according to claim 1, wherein the Go-CDMA matrices allow the introduction of at least some deterministic errors.

20. The method according to claim 2, wherein the Go-CDMA matrices allow the introduction of at least some deterministic errors.

20 21. The method according to claim 3, wherein the Go-CDMA matrices allow the introduction of at least some deterministic errors.

22. A method for decoding a code division multiple access signal based on Go-CDMA codes, comprising:

receiving a signal over a communication channel;

providing majority logic decoding blocks, each block comprising a Go-CDMA matrix;

and

decoding a data message from the signal based on the majority coding blocks.

23. The method according to claim 22, wherein the majority coding blocks comprise a single decoding stage.

24. The method according to claim 22, wherein the majority coding blocks comprise at least two decoding stages.

25. The method according to claim 22, wherein the majority coding blocks comprise at least three decoding stages.

26. The method according to claim 22, wherein the majority coding blocks comprise more than three decoding stages.

27. The method according to claim 22, further comprising decoding a plurality of data messages from the signal based on the majority coding blocks.

28. The method according to claim 27, wherein at least one of the data messages is associated with an active user.

29. The method according to claim 27, wherein at least one of the data messages is associated with a pseudo active user.

30. The method according to claim 27, wherein at least one of the data messages is associated with an inactive user.

31. The method according to claim 24, further comprising a permutation stage between each of at least one adjacent pair of the coding stages.

32. The method according to claim 25, further comprising a permutation stage between each of at least one adjacent pair of the coding stages.

33. The method according to claim 26, further comprising at least one permutation stage between each of at least adjacent pair of the coding stages.

34. The method according to claim 22, wherein the data message comprises data message elements in polar binary format.

35. The method according to claim 22, wherein the data message comprises data message elements in ternary format, or higher order formats.

36. The method according to claim 27, wherein each of the plurality of data messages is based on data received from an intermittent data source.

5 37. The method according to claim 1, wherein the majority coding logic blocks are implemented as a look up table and wherein the coding is performed based on the look up table.

38. The method according to claim 2, wherein the majority coding logic blocks are implemented as a look up table and wherein the coding is performed based on the look up table.

39. The method according to claim 3, wherein the majority coding logic blocks are implemented as a look up table and wherein the coding is performed based on the look up table.

40. The method according to claim 1, wherein the Go-CDMA matrices allow the introduction of at least some deterministic errors.

41. The method according to claim 2, wherein the Go-CDMA matrices allow the introduction of at least some deterministic errors.

20 42. The method according to claim 3, wherein the Go-CDMA matrices allow the introduction of at least some deterministic errors.

ABSTRACT

A class of $n \times \ell$ nonlinear block codes, termed Go-CDMA codes are constructed using column-reduced and row-reduced Hadamard orthogonal matrices, termed Go-CDMA matrices.

Here n, ℓ are positive integers: n chips of user data are transmitted in frames of size $\ell \leq \alpha n$,

5 where α is the frame expansion factor. The codes map n -vectors containing binary message data to binary or multi-level ℓ -vectors for transmission, where $\ell \geq n$. The codes are invertible maps for the binary message data, and when there is no message data in some input vector elements, and noise added between the coding and decoding, there is some error correction. The coding uses integer arithmetic and integer *quantization* operations, preferably certain *sign* operations.

10 Go-CDMA codes may be implemented in CDMA communication systems to improve performance on many measures over conventional CDMA and TDMA systems.

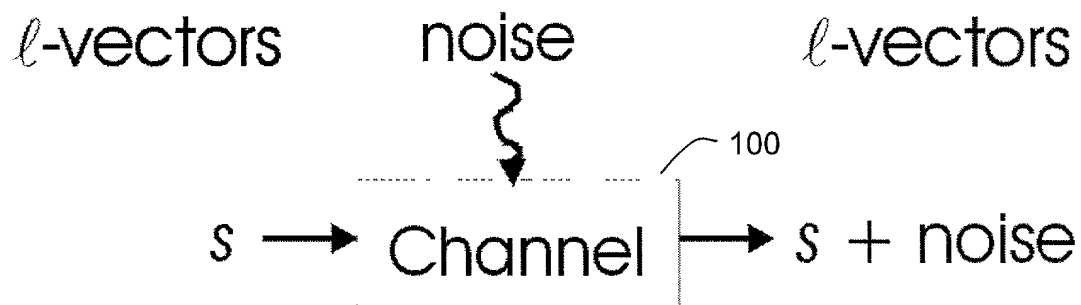


FIGURE 1

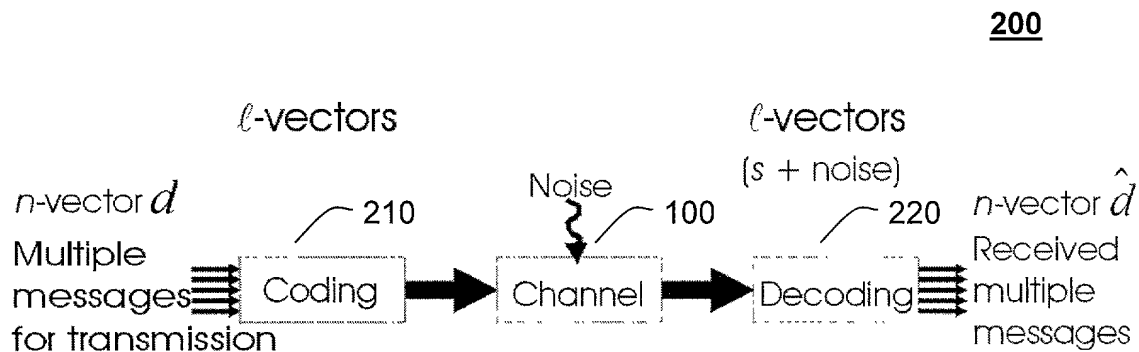


FIGURE 2

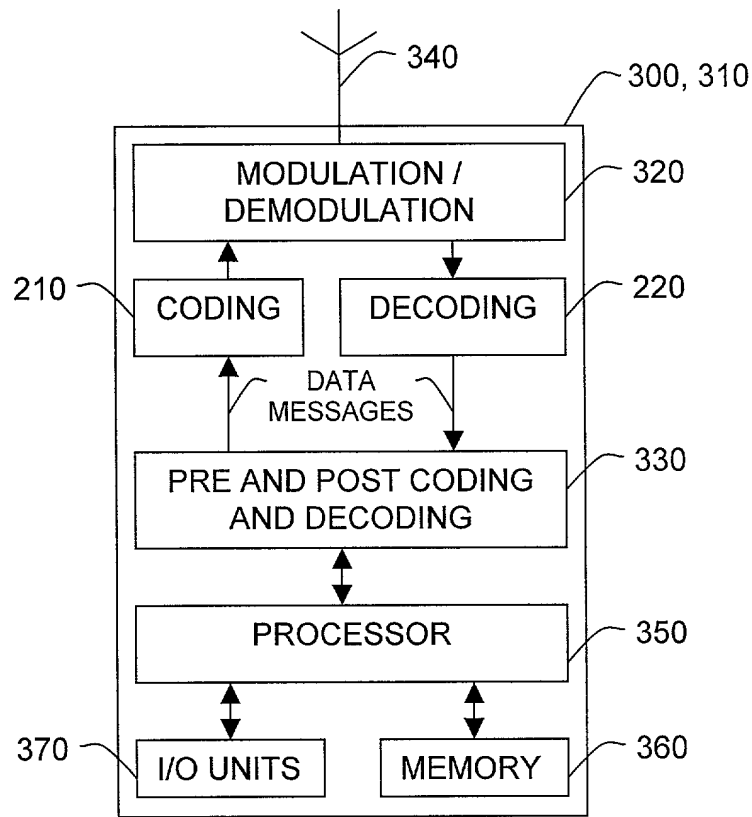


FIGURE 3A

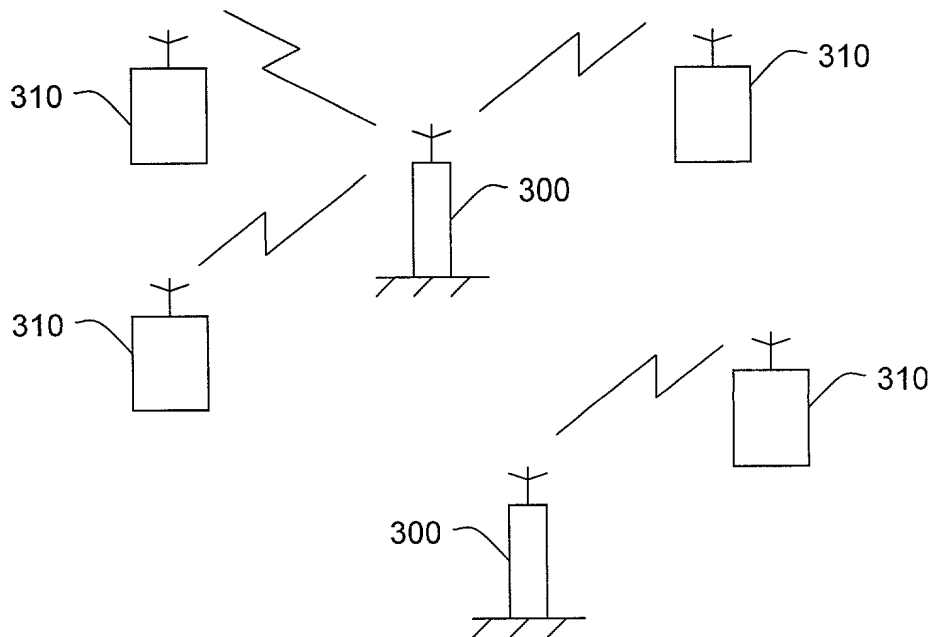


FIGURE 3B

Input

Output

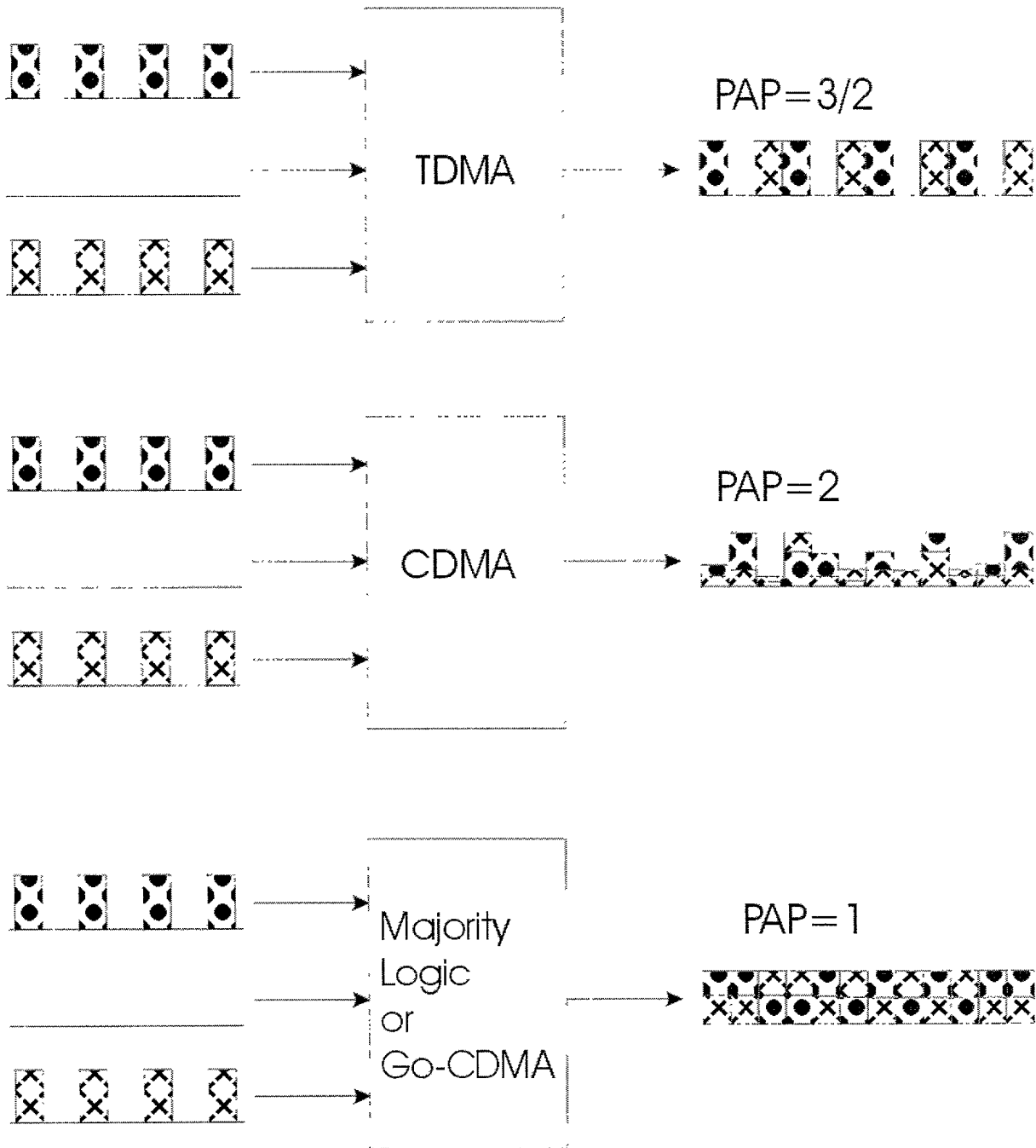


FIGURE 4

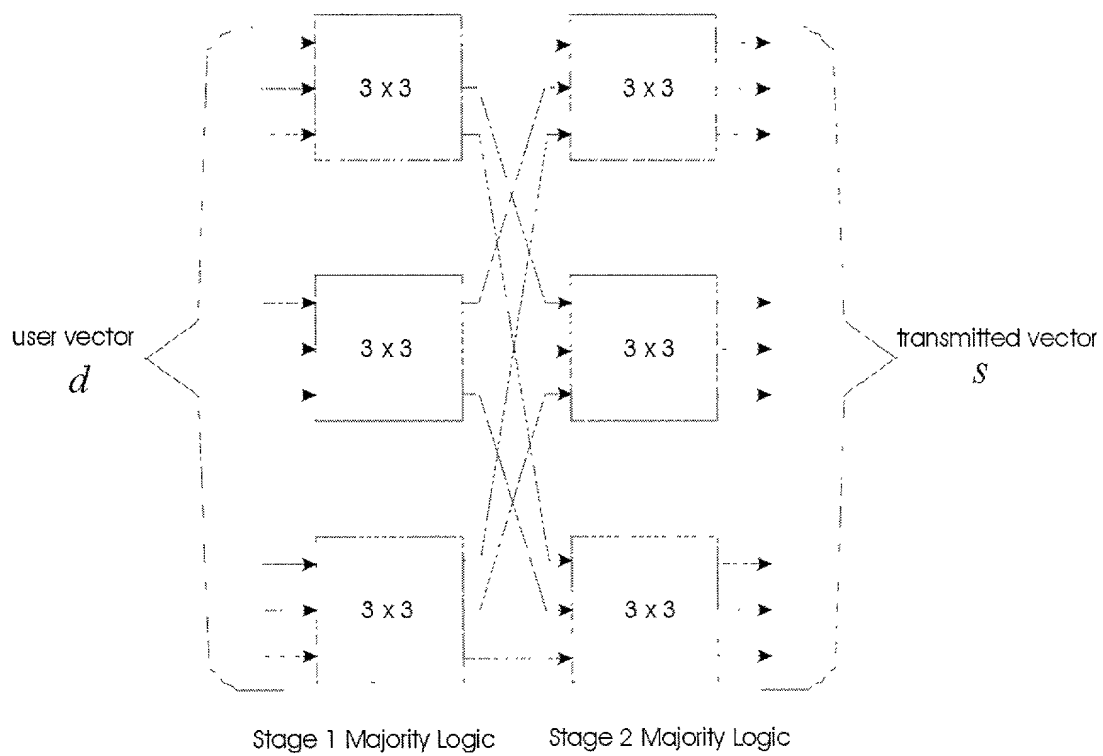


FIGURE 5

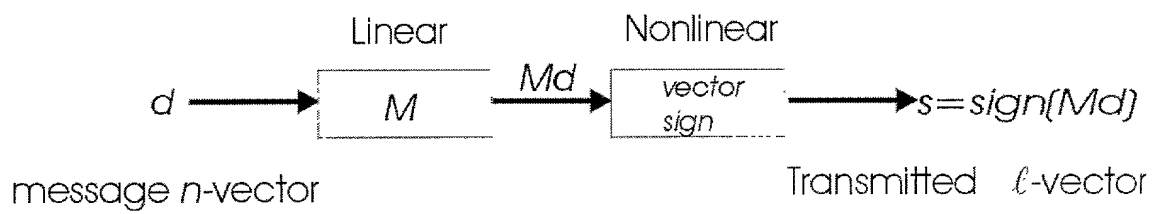


FIGURE 6

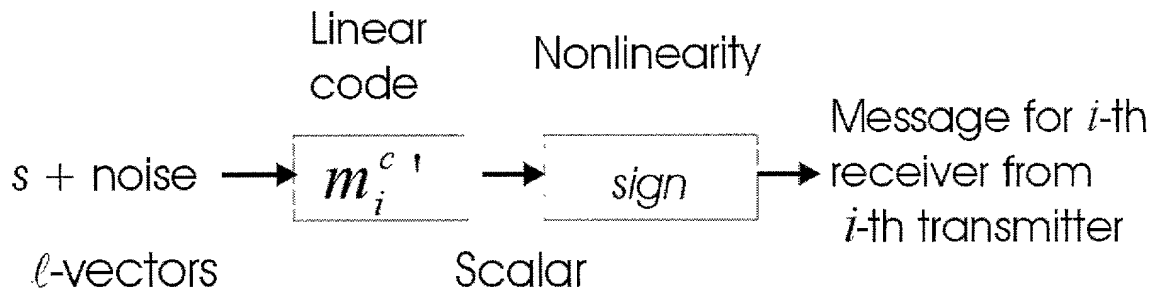


FIGURE 7

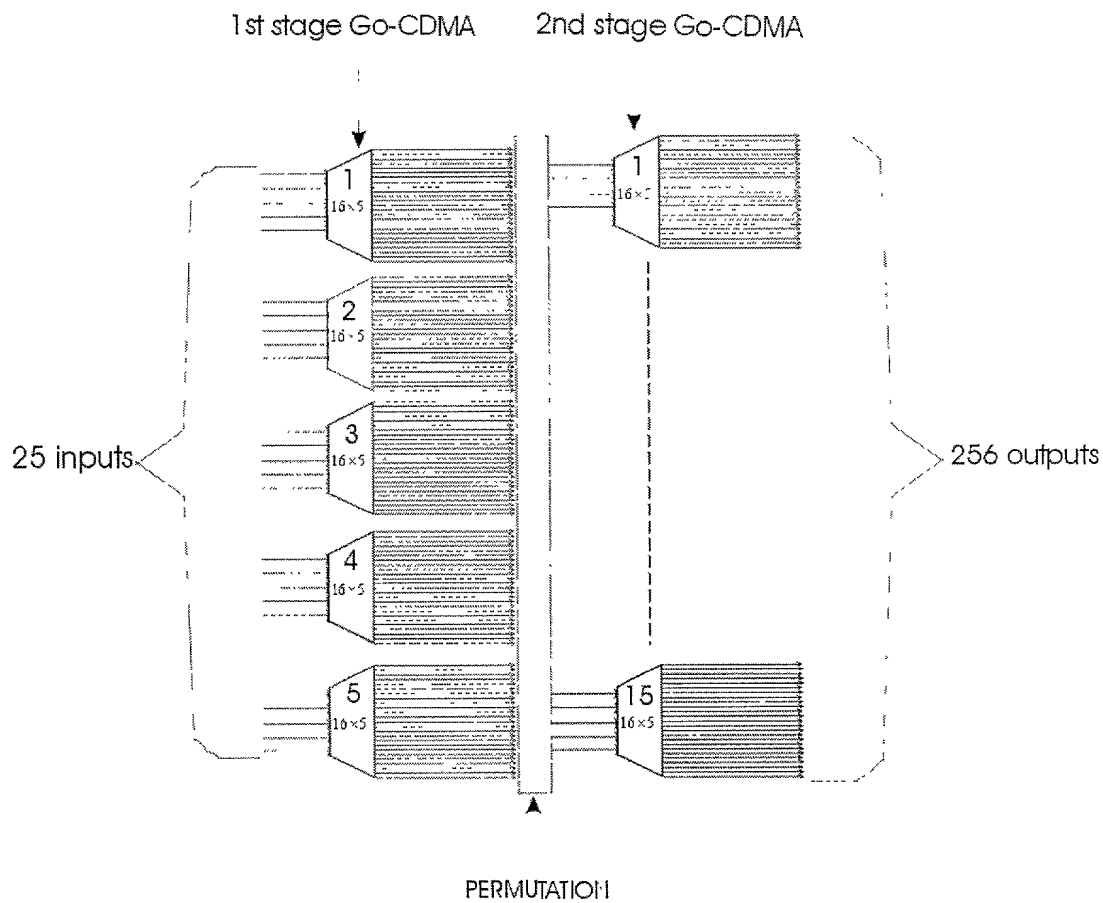


FIGURE 8

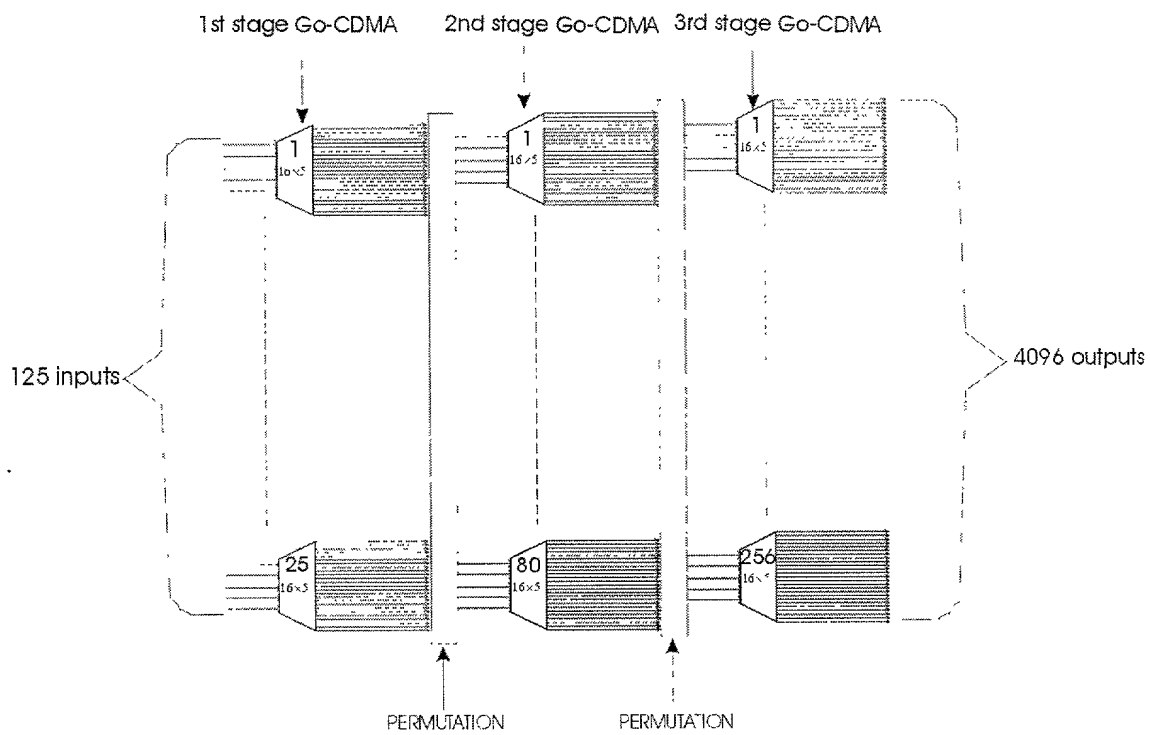


FIGURE 9

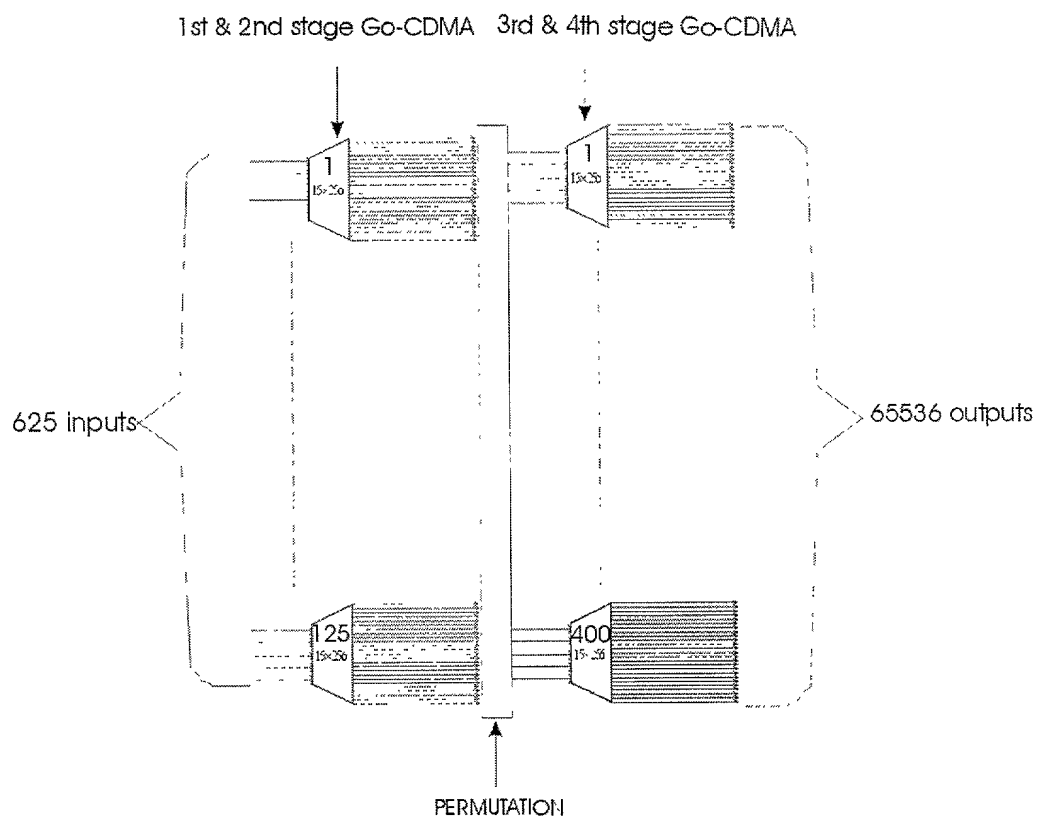


FIGURE 10

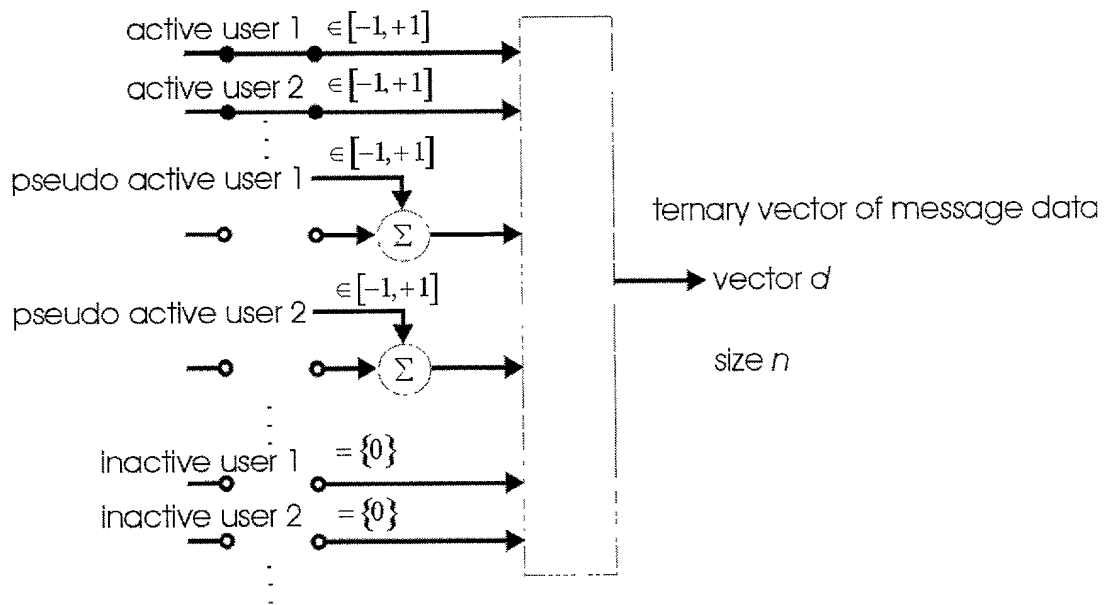


FIGURE 11

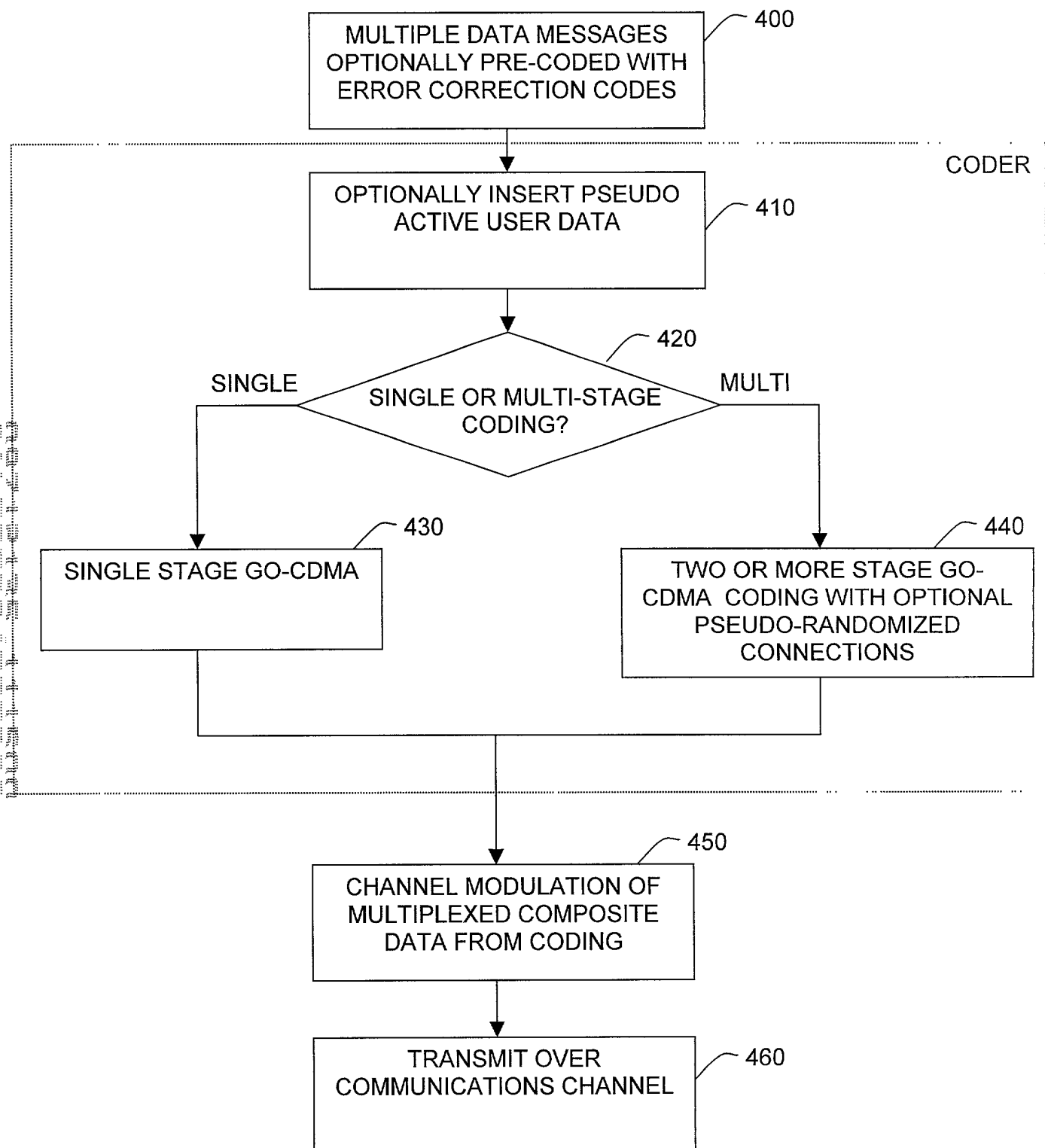


FIGURE 12

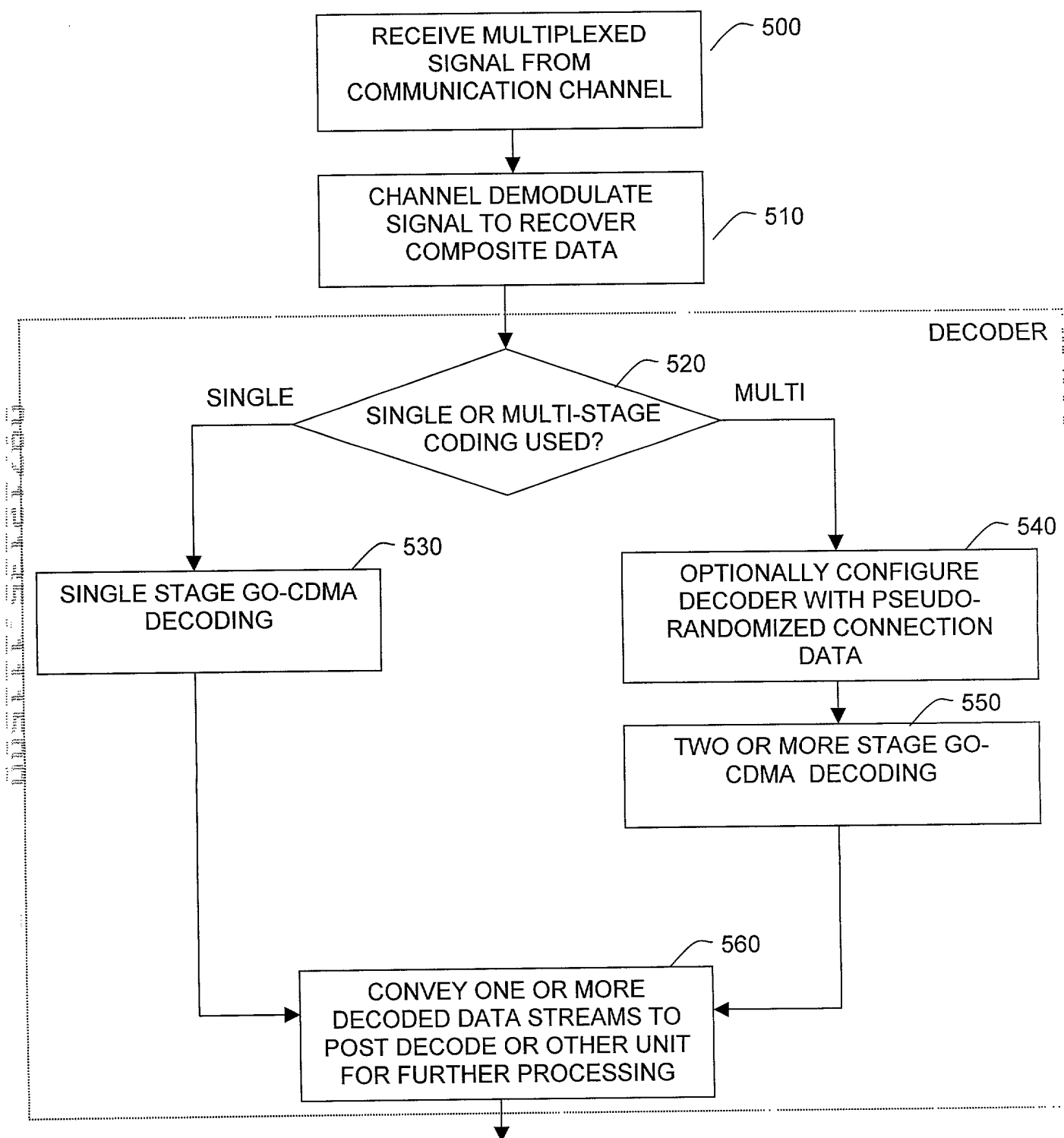


FIGURE 13

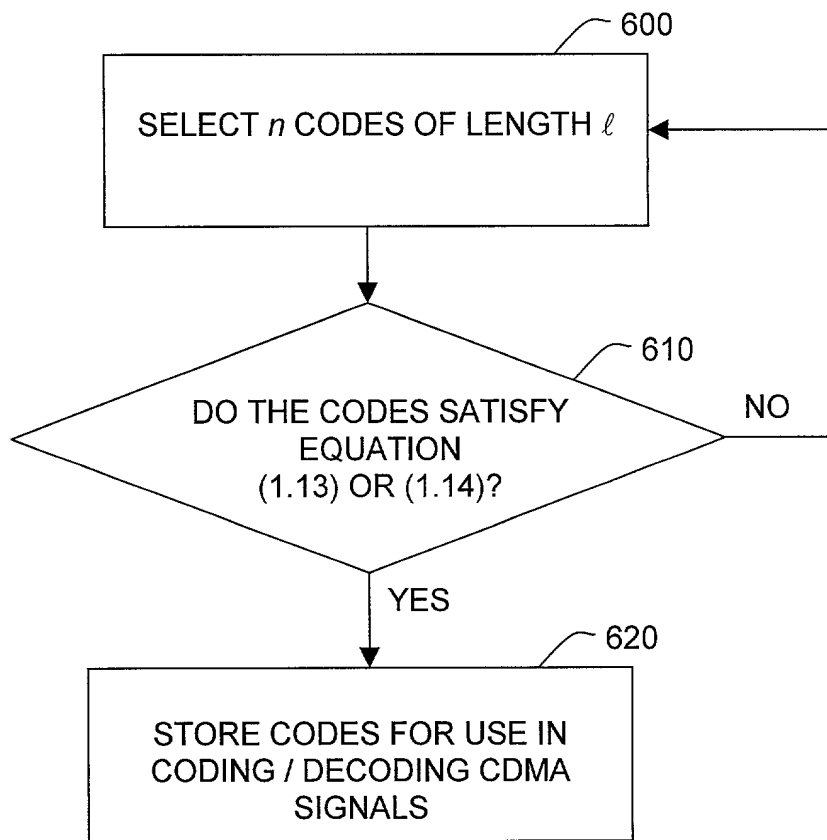


FIGURE 14